

1 **SHARPER BOUNDS FOR PROXIMAL GRADIENT ALGORITHMS**  
2 **WITH ERRORS\***

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4 **Abstract.** We analyse the convergence of the proximal gradient algorithm for convex  
5 composite problems in the presence of gradient and proximal computational inaccura-  
6 cies. We generalize the deterministic analysis to the quasi-Fejér case and quantify the  
7 uncertainty incurred from approximate computing and early termination errors. We  
8 propose new probabilistic tighter bounds that we use to verify a simulated Model Pre-  
9 dictive Control (MPC) with sparse controls problem solved with early termination,  
10 reduced precision and proximal errors. We also show how the probabilistic bounds are  
11 more suitable than the deterministic ones for algorithm verification and more accurate  
12 for application performance guarantees. Under mild statistical assumptions, we also  
13 prove that some cumulative error terms follow a martingale property. And conform-  
14 ing to observations, e.g., in [25], we also show how the acceleration of the algorithm  
15 amplifies the gradient and proximal computational errors.

16 **Key words.** Convex Optimization, Proximal Gradient Descent, Approximate Algorithms

17 **AMS subject classifications.** 49M37, 65K05, 90C25

18 **1. Introduction.** Many problems in science and engineering can be posed as  
19 *composite optimization problems*:

20 (1.1) 
$$\underset{x \in \mathbb{R}^n}{\text{minimize}} \quad f(x) := g(x) + h(x),$$
  
21

22 where the function  $g : \mathbb{R}^n \rightarrow \mathbb{R}$  is real-valued and differentiable, and the function  
23  $h : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\}$  is not necessarily differentiable and is possibly infinite-valued,  
24 enabling the inclusion of hard constraints in (1.1). Examples include various machine  
25 learning frameworks, e.g., logistic regression and support vector machines [11], sparse  
26 regression and inference [23, 15, 16], image processing [1], and discrete optimal control  
27 [17].

28 A popular class of algorithms to solve (1.1) is *proximal gradient methods* [4] which,  
29 in each iteration, take a gradient step using the function  $g$  and, subsequently, evaluate  
30 the proximal operator of the function  $h$  at the resulting point. Such algorithms  
31 have been widely studied under different contexts, and several guarantees have been  
32 established, both in the convex [5, 4, 6, 10, 22] and nonconvex [7, 21] cases. Stochastic  
33 versions of the proximal gradient algorithm have also been proposed and shown to  
34 converge in convex and nonconvex settings, e.g., [2, 29, 20, 24, 12, 30].

35 All of these results, however, assume that computations are performed with near-  
36 infinite precision, which is unrealistic when the computational platform has limitations  
37 in power, precision, or both. Examples include applications that are associated with  
38 sensing and control of autonomous platforms, often using FPGAs or other finite preci-  
39 sion computational hardware. With these applications in mind, we analyze proximal  
40 gradient methods when both the gradient and the proximal operator are computed  
41 approximately at each iteration, and obtain tight performance bounds.

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\*Work supported by UK's EPSRC (EP/T026111/1, EP/S000631/1), and the MOD University Defence Research Collaboration.

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42 While standard proximal gradient methods converge to a solution of (1.1) pro-  
 43 vided the stepsize  $s_k$  is small enough, approximate proximal gradient algorithms re-  
 44 quire, in addition, that the approximation errors  $\epsilon_1^k$  and  $\epsilon_2^k$  satisfy some additional  
 45 convergence criteria, for example, that they converge to zero along the iterations.

46 Our goal is then *to characterize the convergence of the approximate proximal*  
 47 *gradient* to a solution of (1.1). Differently from prior work, we assume not only deter-  
 48 ministic errors, but also probabilistic ones, according to models suited to approximate  
 49 computing.

50 **1.1. Our approach.** In the case of deterministic errors, we get inspiration  
 51 from [4] to derive, using simple arguments, upper bounds on  $f(x^k)$  throughout the  
 52 iterations. The resulting bounds generalize other bounds [25] in the presence of Lip-  
 53 schitz uncertainty and early termination errors under mild assumptions. In the case  
 54 of probabilistic errors, our arguments rely on concentration of measure results for  
 55 martingale sequences and bypass the need to assume that  $\epsilon_1^k$  and  $\epsilon_2^k$  converge to zero.  
 56 The latter yields tighter bounds, and we believe this line of reasoning is novel in the  
 57 analysis of approximate proximal gradient algorithms.

58 **1.2. Applications.** In order to validate our convergence results, we use the pro-  
 59 posed error bounds to analyse the convergence of proximal gradient when applied to  
 60 solve the optimization problem stemming from each time step of Model Predictive  
 61 Control (MPC) [13] with different levels of injected gradient and proximal computa-  
 62 tion errors.

63 **1.3. Contributions.** We summarize our contributions as follows:

- 64 • We establish convergence bounds for the proximal gradient algorithm with  
 65 deterministic and probabilistic errors. Our deterministic bounds generalize  
 66 prior bounds to the quasi-Fejér case where we consider approximate iterations  
 67 and early termination errors and quantify second-order uncertainties. The  
 68 probabilistic bounds tighten the latter under mild conditions.
- 69 • We conduct experiments on a discrete model predictive control problem to  
 70 verify the sharpness of our bounds and compare them with the bounds in [25].  
 71 The models for the errors are inspired by approximate computing techniques  
 72 suited for low-precision machines, such as reduced-precision accelerators on  
 73 FPGA and battery-operated devices, in which algorithms are typically run  
 74 approximately in order to save processing time and/or power.
- 75 • We propose new models for the proximal and gradient errors that satisfy  
 76 martingale properties in accordance with experimental results.

77 **1.4. Organization.** We start by reviewing prior work in Section 2. We then  
 78 describe our approximate computational model, state our assumptions, and present  
 79 the main results in Section 3. The proofs of the main results are included in Section 4,  
 80 and some auxiliary results are relegated to the appendix. Section 5 describes our  
 81 experimental results.

82 **2. Related Work.** One year after the seminal work in [5], it was shown that  
 83 the same nearly optimal rates can still be achieved when the computation of the  
 84 gradients and proximal operators are approximate [25]. This variant is known as the  
 85 *approximate* proximal gradient algorithm. The analysis in [25] requires the errors  
 86  $\epsilon_1^k$  and  $\epsilon_2^k$  to decrease with iterations  $k$  at rates  $O(1/k^{\varsigma+1})$  for the basic proximal  
 87 gradient, and  $O(1/k^{\varsigma+2})$  for the accelerated proximal gradient, for any  $\varsigma > 0$ , in  
 88 order to satisfy the summability assumptions of both error terms. The work in [25]

89 established the following ergodic convergence bound in terms of function values of the  
 90 averaged iterates for the basic approximate proximal gradient (3.7):

$$\begin{aligned}
 & f\left(\frac{1}{k} \sum_{i=1}^k x^i\right) - f(x^*) \leq \frac{L}{2k} \left[ \|x^* - x^0\|_2 + 2A_k + \sqrt{2B_k} \right]^2 \\
 & A_k = \sum_{i=1}^k \left( \frac{\|\epsilon_1^i\|_2}{L} + \sqrt{\frac{2\epsilon_2^i}{L}} \right), \quad B_k = \sum_{i=1}^k \frac{\epsilon_2^i}{L},
 \end{aligned}
 \tag{2.1}$$

92 where  $x^*$  is an optimal solution of (1.1),  $L$  is the Lipschitz constant of the gradient,  
 93 and  $x^0$  is the initialization vector. The same work also analyzed the *approximate*  
 94 *accelerated proximal gradient* and obtained the following convergence result in terms  
 95 of the function values of the iterates,

$$\begin{aligned}
 & f(x^i) - f(x^*) \leq \frac{2L}{(k+1)^2} \left[ \|x^* - x^0\|_2 + 2\tilde{A}_k + \sqrt{2\tilde{B}_k} \right]^2 \\
 & \tilde{A}_k = \sum_{i=1}^k i \left( \frac{\|\epsilon_1^i\|_2}{L} + \sqrt{\frac{2\epsilon_2^i}{L}} \right), \quad \tilde{B}_k = \sum_{i=1}^k \frac{i^2 \epsilon_2^i}{L}.
 \end{aligned}
 \tag{2.2}$$

97 This is the most closely related work to ours; however, our work derives similar, yet  
 98 sharper, convergence bounds. In addition, we derive probabilistic bounds that can  
 99 be estimated before running the algorithm for given bounded proximal and gradient  
 100 errors. Specifically, the constants can be computed from the machine representation  
 101 and software solver tolerances (for the computation of the proximal operator).

102 The work in [3] extended the analysis of [25] to a more general momentum pa-  
 103 rameter selection  $\alpha_k = ((k+a-1)/a)^d$ , where  $d \in [0, 1]$  and  $a > \max(1, (2d)^{\frac{1}{d}})$ ,  
 104 which becomes FISTA [5] when  $d = 1$ . The works in [3, 26] also considered two differ-  
 105 ent types of approximation in the proximal operator computation. For example, [3,  
 106 Proposition 3.3] makes assumptions similar to ours, but establishes different bounds.  
 107 The same paper also suggests slowing down the over-relaxations of FISTA to stabilize  
 108 the algorithm and shows how to obtain a better trade-off between acceleration and  
 109 error amplification by controlling the approximation errors. In contrast, we show that  
 110 the basic approximate proximal gradient algorithm (3.7) converges to a constant pre-  
 111 dictable residual without any assumptions on the gradient error terms (see Theorem  
 112 3). We also show that errors in the accelerated proximal gradient method cause the  
 113 algorithm to eventually diverge as  $O(k)$  in the worst case scenario, but to converge  
 114 sub-optimally, i.e., to a constant error term, using stronger assumptions on the proxi-  
 115 mal error and under a standard suitable choice of the momentum sequence  $\{\beta_k\}$ . We  
 116 also quantify the uncertainties that result from using an inexact optimal reference  
 117 point (motivated by early termination of practical solvers), inexact Fejér monotonic-  
 118 ity (quasi-Fejér monotonicity) and an inexact version of Lipschitz continuity which is  
 119 associated with approximate gradients with the relative error model 3.8.

120 **3. Main Results.** Before stating our convergence guarantees for the approxi-  
 121 mate proximal gradient algorithm, we specify our assumptions and describe the class  
 122 of algorithms that our analysis covers.

123 **3.1. Setup and algorithms.** Recall that we aim to solve convex *composite*  
 124 *optimization problems* with the format of (1.1), repeated here for convenience:

$$\begin{aligned}
 & (3.1) \quad \underset{x \in \mathbb{R}^n}{\text{minimize}} \quad f(x) := g(x) + h(x).
 \end{aligned}$$

127 All of our results assume the following:

128 ASSUMPTION 1 (Assumptions on the problem).

- 129 • The function  $h : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\}$  is closed, proper, and convex.
- 130 • The function  $g : \mathbb{R}^n \rightarrow \mathbb{R}$  is convex and differentiable, and its gradient
- 131  $\nabla g : \mathbb{R}^n \rightarrow \mathbb{R}^n$  is Lipschitz-continuous with constant  $L > 0$ , that is,

$$132 \quad (3.2) \quad \|\nabla g(y) - \nabla g(x)\|_2 \leq L\|y - x\|_2,$$

133 for all  $x, y \in \mathbb{R}^n$ , where  $\|\cdot\|_2$  stands for the standard Euclidean norm.

- 134 • The set of optimal solutions of (3.1) is nonempty:

$$135 \quad (3.3) \quad X^* := \{x \in \mathbb{R}^n : f(x) \leq f(z), \text{ for all } z \in \mathbb{R}^n\} \neq \emptyset.$$

136 The above assumptions are standard in the analysis of proximal gradient algorithms  
137 and are actually required for convergence to an optimal solution from an arbitrary  
138 initialization [4, 6].

139 A consequence of (3.2) that we will often use in our results is that [19, Lem. 1.2.3]

$$140 \quad (3.4) \quad g(y) \leq g(x) + \nabla g(x)^\top (y - x) + \frac{L}{2}\|y - x\|_2^2,$$

141 for any  $x, y \in \mathbb{R}^n$ . Also, as  $h$  is closed, proper, and convex, the function  $z \mapsto h(z) +$   
142  $(1/2)\|z - y\|_2^2$  is coercive, which implies that the approximate set-valued proximal  
143 operator of  $h : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\}$  at  $y \in \mathbb{R}^n$ , defined as

$$144 \quad \text{prox}_h^\epsilon(y) := \left\{x \in \mathbb{R}^n : h(x) + \frac{1}{2}\|x - y\|_2^2 \leq \epsilon + \inf_z h(z) + \frac{1}{2}\|z - y\|_2^2\right\} \neq \emptyset,$$

145 is nonempty for all  $\epsilon \geq 0$ , and  $y \in \mathbb{R}^n$ . When  $\epsilon = 0$ , the proximal operator is computed  
146 exactly, and it is single-valued (a singleton) for closed, proper convex functions

$$147 \quad (3.5) \quad \text{prox}_h(y) := \arg \min_{x \in \mathbb{R}^n} h(x) + \frac{1}{2}\|x - y\|_2^2.$$

148 When  $\epsilon \geq 0$ , this set may contain more than a single element, which results in several  
149 possible instances of the accelerated approximate proximal gradient,

$$150 \quad (3.6) \quad \begin{aligned} y^k &= x^k + \beta_k(x^k - x^{k-1}), \\ x^{k+1} &\in \text{prox}_{s_k h}^{\epsilon_2^k} \left[ y^k - s_k(\nabla g(y^k) + \epsilon_1^k) \right], \end{aligned}$$

151 whenever there exists a  $k$  for which  $\epsilon_2^k > 0$ . However, as we establish bounds on  
152 function values [i.e.,  $f(x^k)$ ], this ambiguity does not affect our results. By setting  
153  $\beta_k = 0$ , (3.6) reduces to the basic approximate proximal gradient scheme, i.e.,

$$154 \quad (3.7) \quad x^{k+1} \in \text{prox}_{s_k h}^{\epsilon_2^k} \left[ x^k - s_k(\nabla g(x^k) + \epsilon_1^k) \right].$$

155 **3.2. Error models and assumptions.** In what follows we consider a relative  
156 error model for the gradient error  $\epsilon_1$ .

157 **ERROR MODEL.** Under this model, each evaluation of the gradient of  $g$  at a point  
158  $x$  is subject to additive noise  $\epsilon_1$  whose magnitude is proportional to the magnitude of  
159 the gradient  $|\nabla g(x)|$ . Specifically, the gradient of  $g$  in (3.1) is approximated by

$$160 \quad (3.8) \quad \nabla g^{\epsilon_1}(x) = \nabla g(x) + \epsilon_1,$$

161 where

$$162 \quad (3.9) \quad |\epsilon_1| \leq \delta |\nabla g(x)|.$$

164  $\delta$  is a positive scalar, and  $|\cdot|$  stands for the vector componentwise absolute value. This  
165 can be used, for example, to model errors in floating-point arithmetic [14].

166 The parameter  $\delta$  is known as the machine precision.

167 For the above error model, our analysis assumes two different scenarios:

- 168 1. The sequences of errors  $\{\epsilon_1^k\}_{k \geq 1}$  and  $\{\epsilon_2^k\}_{k \geq 1}$  are deterministic, or
- 169 2. The sequences of errors  $\{\epsilon_1^k\}_{k \geq 1}$  and  $\{\epsilon_2^k\}_{k \geq 1}$  are random, in which case we  
170 use  $\epsilon_{1\Omega}^k$  and  $\epsilon_{2\Omega}^k$  to denote the respective random vectors/variables of errors at  
171 iteration  $k$ , where  $\Omega$  denotes the sample space of a given probability measure.

172 In scenario 2, the sequences  $\{x^k\}_{k \geq 1}$  and  $\{y^k\}_{k \geq 1}$  become random as well. And we  
173 also use  $x_\Omega^k$  and  $y_\Omega^k$  to denote the respective random vectors at iteration  $k$ . We make  
174 the following assumptions in this case:

175 **ASSUMPTION 2.** In scenario 2, we assume that each random vector  $\epsilon_{1\Omega}^k$ , for  $k \geq 1$ ,  
176 satisfies

$$177 \quad (3.10a) \quad \mathbb{E}[\epsilon_{1\Omega}^k \mid \epsilon_{1\Omega}^1, \dots, \epsilon_{1\Omega}^{k-1}] = \mathbb{E}[\epsilon_{1\Omega}^k] = 0,$$

$$178 \quad (3.10b) \quad \mathbb{P}(|\epsilon_{1\Omega}^k| \leq \delta |\nabla g(x_\Omega^k)|) = 1,$$

$$179 \quad (3.10c) \quad \mathbb{E}[\epsilon_{1\Omega}^{k\top} x_\Omega^k \mid \epsilon_{1\Omega}^1, \dots, \epsilon_{1\Omega}^{k-1}, x_{1\Omega}^1, \dots, x_{1\Omega}^{k-1}] = \mathbb{E}[\epsilon_{1\Omega}^{k\top} x_\Omega^k] = 0,$$

$$180 \quad \text{or } \mathbb{E}[\epsilon_{1\Omega}^k \mid x_\Omega^k] = \mathbb{E}[\epsilon_{1\Omega}^k],$$

182 where  $\delta > 0$  is the machine precision.

183 **ASSUMPTION 3.** Let  $\{x^k\}$  denote the sequence produced by (3.6) or (3.7). We  
184 define the residual error vector at iteration  $k$  as

$$185 \quad (3.11) \quad r^k = x^k - \bar{x}^k,$$

186 where  $\bar{x}^k$  stands for the proximal error-free iterate

$$187 \quad (3.12) \quad \bar{x}^{k+1} := \text{prox}_{sh} \left( x^k - s(\nabla g(x^k) + \epsilon_1^k) \right).$$

188 In scenario 2, we assume

$$189 \quad (3.13a) \quad \mathbb{E}[r_\Omega^k \mid r_\Omega^1, \dots, r_\Omega^{k-1}] = \mathbb{E}[r_\Omega^k] = 0,$$

$$190 \quad (3.13b) \quad \mathbb{E}[r_\Omega^{k\top} x_\Omega^k \mid r_\Omega^1, \dots, r_\Omega^{k-1}, x_{1\Omega}^1, \dots, x_{1\Omega}^{k-1}] = \mathbb{E}[r_\Omega^{k\top} x_\Omega^k] = 0,$$

192 **Remark 3.1.** Lemma 1, stated in the appendix, bounds the norm of the residual  
193 vector  $\|r^k\|_2$  as a function of  $\epsilon_2^k$ ; therefore, bounding  $\epsilon_2^k$  implies bounding  $\|r^k\|_2$ .

194 **3.3. Approximate proximal gradient.** In this section, we consider the ap-  
195 proximate proximal gradient algorithm in (3.7), i.e., without acceleration. We start  
196 by considering deterministic error sequences  $\{\epsilon_1^k\}_{k \geq 1}$  and  $\{\epsilon_2^k\}_{k \geq 1}$ , and then we con-  
197 sider the case in which these sequences are random, as in Assumption 2.

198 **3.3.1. Deterministic errors.** Our first result provides a bound on the ergodic  
199 convergence of the sequence of function values, and decouples the contribution of the  
200 errors in the computation of gradient,  $\epsilon_1^k$ , and in the computation of the proximal  
201 operator,  $\epsilon_2^k$  and  $r^k$ .

202 **THEOREM 1.** Consider problem (3.1) and let Assumption 1 hold. Suppose we  
 203 run the approximate proximal gradient in (3.7) with a fixed stepsize  $s_k := s$  satis-  
 204 fying  $s \leq 1/(L + \delta)$ , for all  $k$ , and under the relative error model in (3.8). Let  
 205 the following stopping criteria hold for  $k \geq k_0$ :  $\epsilon_2^k \leq c_2 \|x^{k+1} - x^k\|_2 \leq c_2 \rho$  and  
 206  $\|\epsilon_1^k\|_2 \leq c_1 \|\nabla g(x^{k+1}) - \nabla g(x^k)\|_2$  where  $\rho$ ,  $c_1$ ,  $c_2$  and  $k_0$  are constants. Then, for  
 207 any  $x^* \in X^*$  and  $k \geq k_0$ , the sequence generated by the approximate proximal gradient  
 208 in (3.7) satisfies

(3.14)

$$209 \quad f\left(\frac{1}{k+1} \sum_{i=0}^k x^{i+1}\right) - f(x^*) \leq \frac{1}{k+1} \left[ \sum_{i=0}^k \epsilon_2^i + \sum_{i=0}^k \left( \|\epsilon_1^i\|_2 + \sqrt{\frac{2\epsilon_2^i}{s}} \right) \|x^* - x^0\|_2 \right. \\ \left. + \frac{1}{2s} \|x^* - x^0\|_2^2 \right] + \frac{1}{k+1} \sum_{i=0}^k \left( \|\epsilon_1^i\|_2 + \sqrt{\frac{2\epsilon_2^i}{s}} \right) \left( \sum_{j=1}^i E^j + iC_\rho \right),$$

210 where  $E^j = \sqrt{\frac{2\epsilon_2^j}{s}} + s \|\epsilon_1^{j-1}\|_2$  and  $C_\rho = \sqrt{2Lc_2\rho} + c_1\rho$ .

211 *Proof.* See Section 4.1.

212 Theorem 1 improves over (2.1) by quantifying the uncertainties associated with the  
 213 Lipschitz and Féjer properties in addition to the ones that stem from proximal and  
 214 gradient errors.

215 *Remark 3.2.* For small perturbations and very small stopping criteria, i.e.,  $\rho \approx 0^1$ ,  
 216 (3.14) can be approximated by

(3.15)

$$217 \quad f\left(\frac{1}{k+1} \sum_{i=0}^k x^{i+1}\right) - f(x^*) \lesssim \frac{1}{k+1} \left[ \sum_{i=0}^k \epsilon_2^i + \sum_{i=0}^k \left( \|\epsilon_1^i\|_2 + \sqrt{\frac{2\epsilon_2^i}{s}} \right) \|x^* - x^0\|_2 \right. \\ \left. + \frac{1}{2s} \|x^* - x^0\|_2^2 \right] - \frac{1}{2s} \sum_{i=0}^k \|r^{i+1}\|_2^2,$$

218 where we have dropped the second order error terms and kept the residual error  
 219 vector explicitly, i.e.,  $-\frac{1}{2s} \sum_{i=0}^k \|r^{i+1}\|_2^2$ , which improves the bound progressively with  
 220 iterations.

221 This result implies that the  $O(1/k)$  convergence rate is still guaranteed with  
 222 weaker summability assumptions on  $\{\epsilon_2^k\}_{k \geq 1}$  and  $\{\|\epsilon_1^k\|_2\}_{k \geq 1}$ . For instance, con-  
 223 sider the case where both proximal and gradient errors decrease as  $O(1/k)$  (i.e., non-  
 224 summable). Then Theorem 1 yields an overall convergence rate of  $O(\log k/k)$  which  
 225 is less conservative than what would have been obtained from (2.1), i.e.,  $O(\log^2 k/k)$ .  
 226 Consequently, as a necessary condition for convergence, we only require the partial  
 227 sums  $\sum_{i=1}^k \epsilon_2^i$  and  $\sum_{i=1}^k \|\epsilon_1^i\|_2$  to be in  $o(k)$  as compared to the stronger condition  
 228  $o(\sqrt{k})$  that is implied by (2.1). If we set both errors to zero for all  $k \geq 1$ , we recover  
 229 the error-free optimal upper bound  $\frac{1}{2sk} \|x^* - x^0\|_2^2$  [4].

230 **3.3.2. Random errors.** Let us now consider the case in which  $\epsilon_1^k$ ,  $\epsilon_2^k$  and there-  
 231 fore  $x^k$ , are random, and let  $\epsilon_{1\Omega}^k$ ,  $\epsilon_{2\Omega}^k$  and  $x_\Omega^k$  be the corresponding random vari-  
 232 ables/vectors.

<sup>1</sup> $C_\rho = 0$  if the optimum  $x^*$  is reached.

233 **THEOREM 2 (Random errors).** Consider problem (3.1) and let Assumption 1  
 234 hold. Assume that the gradient error  $\{\epsilon_{1\Omega}^k\}_{k \geq 1}$  and residual proximal error  $\{r_\Omega^k\}_{k \geq 1}$   
 235 sequences satisfy Assumptions 2, 3 and  $\mathbb{P}(\epsilon_{2\Omega}^k \leq \varepsilon_0) = 1$ , for all  $k > 0$ , and for some  
 236  $\varepsilon_0 \in \mathbb{R}$ . Let  $\{x_\Omega^i\}$  denote a sequence generated by the approximate proximal gradient  
 237 algorithm in (3.7) with constant stepsize  $s_k = s \leq 1/(L + \delta)$ , for all  $k$ . Assume that  
 238 there is a positive scalar  $D_x > 0$  such that  $\|x_\Omega^k - x_\Omega^*\|_2^2 \leq D_x \|x_\Omega^0 - x_\Omega^*\|_2^2$  holds with  
 239 probability  $p$ , for all  $k$ . Then, for any  $\gamma > 0$ ,

(3.16)

$$\begin{aligned}
 240 \quad f\left(\frac{1}{k} \sum_{i=1}^k x_\Omega^i\right) - f(x^*) &\leq \frac{1}{k} \sum_{i=1}^k \epsilon_{2\Omega}^i + \frac{\gamma}{\sqrt{k}} \left( \sqrt{n} M_{\nabla g} |\delta| + \sqrt{\frac{2\varepsilon_0}{s}} \right) D_x \|x^* - x^0\|_2 \\
 &+ \frac{D_x^2}{2sk} \|x^* - x^0\|_2^2,
 \end{aligned}$$

241 with probability at least  $p^k (1 - 2 \exp(-\frac{\gamma^2}{2}))$ , where  $x^*$  is any solution of (3.1),  $M_{\nabla g} =$

$$242 \sup_{i \in \mathbb{N}_+} \left\{ \|\nabla g(x^i)\|_\infty \right\}.$$

243 *Proof.* See Section 4.2

244 For large scale problems,<sup>2</sup> we typically have  $n \gg \frac{1}{s} \geq L$ ; therefore, we obtain the  
 245 following approximated bound

(3.17)

$$246 \quad f\left(\frac{1}{k} \sum_{i=1}^k x_\Omega^i\right) - f(x^*) \approx \frac{1}{k} \sum_{i=1}^k \epsilon_{2\Omega}^i + \gamma M_{\nabla g} D_x \sqrt{\frac{n}{k}} |\delta| \|x^* - x^0\|_2 + \frac{D_x^2}{2sk} \|x^* - x^0\|_2^2,$$

247 with approximately the same probability. In the absence of computational errors,  
 248 (3.16) reduces to the deterministic noise-free convergence bound for  $D_x = 1$ , i.e.,  
 249  $\frac{1}{2sk} \|x^* - x^0\|_2^2$ .

250 The following result applies if we assume statistical stationarity<sup>3</sup> of proximal  
 251 errors.

252 **THEOREM 3 (Random stationary errors).** Consider problem (3.1), let As-  
 253 sumptions 1 hold and assume that the rounding error  $\{\epsilon_{1\Omega}^k\}_{k \geq 1}$  and residual error  
 254  $\{r_\Omega^k\}_{k \geq 1}$  sequences satisfy Assumptions 2, 3 and that the proximal computation error  
 255 is upper bounded, i.e.  $\mathbb{P}(\epsilon_{2\Omega}^k \leq \varepsilon_0) = 1$  for all  $k \geq 1$  and stationary with constant mean  
 256  $\mathbb{E}[\epsilon_{2\Omega}]$ . Let  $\{x_\Omega^i\}$  denote a sequence generated by the approximate proximal gradient  
 257 algorithm in (3.7) with constant stepsize  $s_k = s \leq 1/(L + \delta)$ , for all  $k$ . Assume that  
 258 there is a positive scalar  $D_x > 0$  such that  $\|x_\Omega^k - x_\Omega^*\|_2^2 \leq D_x^2 \|x_\Omega^0 - x_\Omega^*\|_2^2$  holds with  
 259 probability  $p$ , for all  $k$ . Then, for any  $\gamma > 0$ ,

$$\begin{aligned}
 260 \quad (3.18) \quad f\left(\frac{1}{k} \sum_{i=1}^k x_\Omega^i\right) - f(x^*) &\leq \mathbb{E}(\epsilon_{2\Omega}) + \frac{\gamma}{\sqrt{k}} \left( \frac{\varepsilon_0}{2} + \sqrt{n} M_{\nabla g} D_x |\delta| \|x^* - x^0\|_2 \right) \\
 &+ \frac{D_x^2}{2sk} \|x^* - x^0\|_2^2,
 \end{aligned}$$

<sup>2</sup>And for same levels of error magnitudes  $\delta$  and  $\varepsilon_0$ .

<sup>3</sup>Whose ensemble mean and variance are time-invariant.

261 with probability at least  $p^k(1 - 4\exp(-\frac{\gamma^2}{2}))$ , where  $x^*$  is any solution of (3.1),  $M_{\nabla g} =$   
 262  $\sup_{i \in \mathbb{N}_+} \left\{ \|\nabla g(x^i)\|_\infty \right\}$ .

263 *Proof.* See Section 4.3

264 *Remark 3.3.*  $D_x$  could be taken as large as to satisfy  $\|x_\Omega^k - x_\Omega^*\|_2^2 \leq D_x^2 \|x_\Omega^0 - x_\Omega^*\|_2^2$   
 265 almost surely, i.e., with probability 1.

266 Once again, if both errors are forced to zero in (3.18) then the optimal convergence  
 267 rate is obtained as in Theorem 1 and Theorem 2. (3.18) also implies that we obtain  
 268 a worst case convergence rate of  $O(1)$ , i.e., convergence up to a predicted constant  
 269 residual  $\mathbb{E}[\epsilon_{2\Omega}]$ .

### 270 3.4. Accelerated Approximate PG.

271 **3.4.1. Deterministic errors.** We now analyze the effect of computational in-  
 272 accuracy on the approximate accelerated PG. In what follows, we establish upper  
 273 bounds on the convergence of the accelerated PG in the presence of deterministic  
 274 errors in the computation of the gradient as well as in the proximal operation step.

275 **THEOREM 4 (Accelerated with deterministic errors).** *Consider problem*  
 276 *(3.1) and let Assumption 1 hold. Suppose we run the approximate accelerated proxi-*  
 277 *mal gradient in (3.6) with a fixed stepsize  $s_k := s$  satisfying  $s \leq 1/(L + \delta)$ , for all  $k$ ,*  
 278 *and under the relative error model in (3.8). Let the following stopping criteria*  
 279 *hold for  $k \geq k_0$ :  $\epsilon_2^k \leq c_2 \|x^{k+1} - x^k\|_2 \leq c_2 \rho$  and  $\|\epsilon_1^k\|_2 \leq c_1 \|\nabla g(x^{k+1}) - \nabla g(x^k)\|_2$*   
 280 *where  $\rho$ ,  $c_1$ ,  $c_2$  and  $k_0$  are constants. Assume we have summable iterative displace-*  
 281 *ments  $\|x^k - x^{k-1}\|_2$ . Let the momentum sequence  $\beta_k = (\alpha_{k-1} - 1)/\alpha_k$  be designed*  
 282 *such that  $\alpha_k$  satisfies the following:*

- 283 •  $\alpha_k \geq 1 \quad \forall \quad k > 0$  and  $\alpha_0 = 1$
- 284 •  $\alpha_k^2 - \alpha_k = \alpha_{k-1}$
- 285 •  $\{\alpha_k\}_{k=0}^\infty$  is an increasing sequence and proportional to  $k$  ( $O(k)$ )

286 *Then, for any  $x^* \in X^*$  and  $k \geq k_0$ , the sequence generated by the approximate*  
 287 *accelerated proximal gradient in (3.6) satisfies*

$$\begin{aligned}
 (3.19) \quad f(x^{k+1}) - f(x^*) &\leq \frac{1}{\alpha_k^2} \left[ \sum_{i=0}^k \alpha_i^2 \epsilon_2^i + \sum_{i=0}^k \alpha_i \|x^0 - x^*\|_2 \left( \|\epsilon_1^i\|_2 + \sqrt{\frac{2\epsilon_2^i}{s}} \right) \right. \\
 &\quad \left. + \frac{1}{2s} \|x^0 - x^*\|_2^2 \right] + \frac{1}{\alpha_k^2} \sum_{i=0}^k \alpha_i \left( \|\epsilon_1^i\|_2 + \sqrt{\frac{2\epsilon_2^i}{s}} \right) \sum_{j=1}^i \alpha_j (E^j + C_\rho),
 \end{aligned}$$

289 where  $x^*$  is any solution of (3.1),  $E^j = \sqrt{\frac{2\epsilon_2^j}{s}} + s \|\epsilon_1^{j-1}\|_2$  and  $C_\rho = \sqrt{2Lc_2\rho} + c_1\rho$ ,  
 290 and  $C_\rho = \sqrt{2Lc_2\rho} + c_1\rho$ .

291 *Proof.* See Section 4.4

292 *Remark 3.4.* Ignoring second order error terms (for small square summable per-  
 293 turbations and very small suboptimality stopping criterion, i.e.,  $\rho \approx 0$ ), (3.19) can be



294 approximated by

$$295 \quad (3.20) \quad f(x^{k+1}) - f(x^*) \lesssim \frac{1}{\alpha_k^2} \left[ \sum_{i=0}^k \alpha_i^2 \epsilon_2^i + \sum_{i=0}^k \alpha_i \left( \|\epsilon_1^i\|_2 + \sqrt{\frac{2\epsilon_2^i}{s}} \right) \|x^0 - x^*\|_2 \right. \\ \left. + \frac{1}{2s} \|x^0 - x^*\|_2^2 \right].$$

296 Notice that if we trivially choose  $\beta_k = 0$  we recover back the nonaccelerated basic  
297 scheme. In the noise-free case, (3.19) reduces to  $\frac{1}{2s\alpha_k^2} \|x^* - x^0\|_2^2$ , which coincides  
298 with the convergence rate of the accelerated proximal gradient algorithm [4, Thm.  
299 10.34], i.e.,  $O(1/k^2)$  if  $\alpha_k$  is in the order of  $O(k)$ .

300 **3.4.2. Random errors.** The following result gives an estimate of the conver-  
301 gence rate when both errors are stochastic and bounded following a probabilistic  
302 analysis approach.

303 **THEOREM 5 (Accelerated with random errors).** *Consider problem (3.1)*  
304 *and let Assumption 1 hold. Suppose that the rounding error  $\{\epsilon_{1\Omega}^k\}_{k \geq 1}$  and residual*  
305 *error  $\{r_\Omega^k\}_{k \geq 1}$  sequences satisfy Assumptions 2 and 3, respectively. Let the norm of*  
306 *the iterative difference  $\|x_\Omega^k - x_\Omega^{k-1}\|_2$  be summable. Define a new sequence  $u_\Omega^k :=$*   
307  *$x^* - x_\Omega^k + (1 - \alpha_{k-1})(x_\Omega^k - x_\Omega^{k-1})$ . Assume that there is a positive scalar  $D_u > 0$  such*  
308 *that  $\|u_\Omega^i\|_2^2 \leq D_u^2 \|x^0 - x^*\|_2^2$  holds with probability  $p$ . Let  $\varepsilon_0$  be an upper bound on the*  
309 *proximal error, i.e.,  $\epsilon_{2\Omega}^k \leq \varepsilon_0$  for all  $k$ . Then, for all  $\gamma > 0$ , the sequence generated*  
310 *by the approximate APG in (3.6) with constant stepsize  $s_k := s \leq 1/(L + \delta)$ , for all*  
311  *$k$ , under error models (3.10) and (3.13), and with the following choices:*

- 312 •  $\beta_k = \frac{\alpha_{k-1} - 1}{\alpha_k}$
- 313 •  $\alpha_k \geq 1 \quad \forall k > 0$  and  $\alpha_0 = 1$
- 314 •  $\alpha_k^2 - \alpha_k = \alpha_{k-1}$
- 315 •  $\{\alpha_k\}_{k=0}^\infty$  increases as  $o(k)$

316 satisfies

$$317 \quad (3.21) \quad f(x_\Omega^{k+1}) - f(x^*) \leq \frac{1}{\alpha_k^2} \left[ S_{\epsilon_{2\Omega}} + S_{r_\Omega} + S_{\epsilon_{1\Omega}} + \frac{1}{2s} \|x^* - x^0\|_2^2 \right],$$

318 where

$$319 \quad (3.22) \quad S_{\epsilon_{2\Omega}} = \varepsilon_0 \sum_{i=0}^k i^2 + \frac{\gamma}{2} \sqrt{\sum_{i=1}^k i^4 (\epsilon_{2\Omega}^i)^2},$$

$$320 \quad (3.23) \quad S_{\epsilon_{1\Omega}} = \gamma |\delta| M_{\nabla g} D_u^2 \|x^0 - x^*\|_2^2 \sqrt{n \sum_{i=1}^k i^2},$$

$$321 \quad (3.24) \quad S_{r_\Omega} = \gamma D_u^2 \|x^0 - x^*\|_2^2 \sqrt{\frac{2}{s} \sum_{i=1}^k i^2 \epsilon_2^i}$$

323 with probability at least  $p^k (1 - 4 \exp(-\gamma^2/2))$ , where  $x^*$  is any solution of (3.1),

324  $M_{\nabla g} = \sup_{i \in \mathbb{N}_+} \left\{ \|\nabla g(x^i)\|_\infty \right\}$ , and  $\mathbb{E}[\cdot]$  stands for the expectation operator.

325 *Proof.* See Section 4.5

326 *Remark 3.5.*  $D_u$  could be taken as large as to satisfy  $\|u_\Omega^i\|_2^2 \leq D_u^2 \|x^0 - x^*\|_2^2$   
 327 with probability 1.

328 The following corollary results from the substitution of partial sums by their cor-  
 329 responding closed forms and using the worst case upper bound  $\varepsilon_0$  on  $\varepsilon_{2\Omega}^i$  for all  
 330  $i = 1, \dots, k$ .

331 **COROLLARY 5.1 (Accelerated with random errors).** *Consider problem (3.1)*  
 332 *and let the assumptions of Theorem 5 hold. Define a new sequence  $w_\Omega^k := x^* - x_\Omega^k +$*   
 333  *$(1 - \alpha_{k-1})(x_\Omega^k - x_\Omega^{k-1})$ . Assume that there is a positive scalar  $D_u > 0$  such that*  
 334  *$\|u_\Omega^i\|_2^2 \leq D_u^2 \|x^0 - x^*\|_2^2$  holds with probability  $p$ . Then we have, for all  $k$ . Let  $\varepsilon_0$  be*  
 335 *an upper bound on the proximal error, i.e.,  $\varepsilon_{2\Omega}^k \leq \varepsilon_0$  for all  $k$ . Then we have, for all*  
 336  $k$ ,

$$337 \quad (3.25) \quad f(x_\Omega^{k+1}) - f(x^*) \leq \frac{1}{\alpha_k^2} \left[ \bar{S}_{\varepsilon_{2\Omega}} + \bar{S}_{r_\Omega} + \bar{S}_{\varepsilon_{1\Omega}} + \frac{1}{2s} \|x^* - x^0\|_2^2 \right],$$

338 where

$$339 \quad (3.26) \quad \bar{S}_{\varepsilon_{2\Omega}} = \varepsilon_0 \frac{k(k+1)(2k+1)}{6} + \frac{\gamma}{2} \varepsilon_0 \sqrt{\frac{k(k+1)(2k+1)(3k^2+3k-1)}{30}},$$

$$340 \quad (3.27) \quad \bar{S}_{\varepsilon_{1\Omega}} = \gamma |\delta| D_u M_{\nabla g} \|x^0 - x^*\|_2 \sqrt{\frac{nk(k+1)(2k+1)}{6}},$$

$$341 \quad (3.28) \quad \bar{S}_{r_\Omega} = \gamma D_u \|x^0 - x^*\|_2 \sqrt{\frac{2s\varepsilon_0 k(k+1)(2k+1)}{6}}.$$

343 with probability at least  $1 - 4 \exp(-\gamma^2/2)$ , where  $x^*$  is any solution of (3.1),  $M_{\nabla g} =$   
 344  $\sup_{i \in \mathbb{N}_+} \left\{ \|\nabla g(x^i)\|_\infty \right\}$ .

345 *Proof.* Substituting

$$346 \quad (3.29) \quad \sum_{i=1}^k i^2 = \frac{k(k+1)(2k+1)}{6},$$

347 and substituting

$$348 \quad (3.30) \quad \sum_{i=1}^k i^4 = \frac{k(k+1)(2k+1)(3k^2+3k-1)}{30},$$

349 and using  $\|u_\Omega^i\|_2 \leq D_u \|x^0 - x^*\|_2$ ,  $\|\varepsilon_{1\Omega}^i\|_2 \leq |\delta| M_{\nabla g} \sqrt{n}$  in Theorem 5 completes the  
 350 proof.  $\square$

351 In the absence of errors, both probabilistic and deterministic analyses lead to the op-  
 352 timal convergence rate of  $O(1/k^2)$  for the accelerated scheme (3.19)-(3.21). However,  
 353 as stated previously in Theorem 5, under the influence of computational inaccura-  
 354 cies and due to error amplification, acceleration has a counter-effect in the Nesterov's  
 355 sense [18] and the method becomes more sensitive to gradient and proximal errors  
 356 whenever we want to speed up the algorithm.

357 Although computational errors are deterministic in nature [14], probabilistic re-  
 358 sults such as (3.21) give us practical convergence bounds when errors cannot be mea-  
 359 sured or are undetectable but with known upper bounds. If the ensemble mean  $\mathbb{E}[\varepsilon_{2\Omega}^k]$

is constant for all  $k \geq 1$  in (3.21), i.e., the error sequence  $\{\epsilon_{2\Omega}^k\}$  is stationary, then (3.21) becomes totally independent from the instantaneous running errors  $\epsilon_{1\Omega}^k, \epsilon_{2\Omega}^k$  as well as from the running iterates  $x_\Omega^k$  and would be only determined by the machine precision  $\delta$ , the tolerance  $\mathbb{E}[\epsilon_{2\Omega}]$  and the given probability parameter  $\gamma$ . The factor  $\alpha_k$  is designed to be proportional to the iteration counter  $o(k)$ .

Although boundedness of the gradient error is sufficient for the gradient error term  $S_{\epsilon_{1\Omega}}$  to asymptotically vanish, the algorithm fails to converge without the summability of the proximal error term  $\{\alpha_k^2 \mathbb{E}(\epsilon_{2\Omega}^k)\}$ .

#### 4. Proofs.

**4.1. Proof of Theorem 1.** Recall the definition of  $\epsilon$ -suboptimal proximal operator in (3.1):

$$(4.1) \quad \text{prox}_u^\epsilon(y) := \left\{ x \in \mathbb{R}^n : u(x) + \frac{1}{2} \|x - y\|_2^2 \leq \epsilon + \inf_z u(z) + \frac{1}{2} \|z - y\|_2^2 \right\}.$$

Because this is a set, the point  $x^{k+1}$  in approximate proximal gradient (3.7) is not defined uniquely. To bound the effect of the error  $\epsilon_2^k$ , we will therefore compute its difference with respect to the case where  $\epsilon_2^k = 0$ , as measured by a function that we will define shortly. Recall that  $\bar{x}^{k+1}$  is the noiseless computation of the proximal operator in (3.7) at  $x^k$  with constant stepsize  $s$ :

$$(4.2) \quad \bar{x}^{k+1} := \text{prox}_{sh} \left[ x^k - s(\nabla g(x^k) + \epsilon_1^k) \right],$$

$$(4.3) \quad = \text{prox}_{sh} \left[ x^k - s\nabla^{\epsilon_1^k} g(x^k) \right]$$

$$(4.4) \quad = \arg \min_x g(x^k) + \nabla^{\epsilon_1^k} g(x^k)^\top (x - x^k) + \frac{1}{2s} \|x - x^k\|_2^2 + h(x)$$

$$(4.5) \quad := \arg \min_x G(x, x^k).$$

From (4.2) to (4.3), we used  $\nabla^{\epsilon_1^k} g(x^k) := \nabla g(x^k) + \epsilon_1^k$  as the inexact gradient of  $g$  at  $x^k$ . From (4.3) to (4.4), we developed the squared  $\ell_2$ -norm term in the definition of the proximal operator [cf. (3.5)] and added  $g(x^k)$  to the objective function. Finally, from (4.4) to (4.5), we defined

$$(4.6) \quad G(x, x^k) := g(x^k) + \nabla^{\epsilon_1^k} g(x^k)^\top (x - x^k) + \frac{1}{2s} \|x - x^k\|_2^2 + h(x).$$

As  $h$  is convex [cf. Assumption 1], the quadratic term in (4.6) makes the function  $G(\cdot, x^k)$  strongly convex with parameter  $1/s$  [4].

Recall that  $\bar{x}^{k+1}$  is the optimal solution of (4.5) and that  $x^{k+1}$  is the actual, noisy iterate in (3.7). Therefore, according to (3.7) and to the definition of the  $\epsilon$ -suboptimal proximal operator in (4.1),

$$(4.7) \quad h(x^{k+1}) + \frac{1}{2s} \left\| x^{k+1} - x^k + s\nabla^{\epsilon_1^k} g(x^k) \right\|_2^2 \leq \epsilon_2^k + h(\bar{x}^{k+1})$$

$$+ \frac{1}{2s} \left\| \bar{x}^{k+1} - x^k + s\nabla^{\epsilon_1^k} g(x^k) \right\|_2^2$$

$$(4.8) \quad \iff h(x^{k+1}) + \frac{1}{2s} \|x^{k+1} - x^k\|_2^2 + \nabla^{\epsilon_1^k} g(x^k)^\top (x^{k+1} - x^k) \leq$$

$$\epsilon_2^k + h(\bar{x}^{k+1}) + \frac{1}{2s} \|\bar{x}^{k+1} - x^k\|_2^2 + \nabla^{\epsilon_1^k} g(x^k)^\top (\bar{x}^{k+1} - x^k)$$

$$(4.9) \quad \iff G(x^{k+1}, x^k) - G(\bar{x}^{k+1}, x^k) \leq \epsilon_2^k.$$

398 From (4.7) to (4.8), we developed the squared-norm terms and cancelled the common  
 399 term. From (4.8) to (4.9), we added the constant  $g(x^k) - \frac{s}{2} \|\nabla g(x^k)\|_2^2$  to both sides  
 400 and used the definition (4.6). Notice that (4.9) bounds the distance between  $x^{k+1}$   
 401 and  $\bar{x}^{k+1}$  as measured by  $G(\cdot, x^k)$ .

402 Because  $G(\cdot, x^k)$  is strongly convex, [4, Theorem. 5.25] establishes that

$$403 \quad (4.10) \quad G(x, x^k) - G(\bar{x}^{k+1}, x^k) \geq \frac{1}{2s} \|x - \bar{x}^{k+1}\|_2^2,$$

404 for any  $x \in \mathbb{R}^n$ . In particular, it holds for any optimal solution  $x^*$  of (3.1).

405 Thus, subtracting (4.10) with  $x = x^*$  from (4.9) yields

$$406 \quad (4.11) \quad G(x^{k+1}, x^k) - G(x^*, x^k) \leq \epsilon_2^k - \frac{1}{2s} \|x^* - \bar{x}^{k+1}\|_2^2$$

$$407 \quad (4.12) \iff g(x^k) + \nabla^{\epsilon_1^k} g(x^k)^\top (x^{k+1} - x^k) + \frac{1}{2s} \|x^{k+1} - x^k\|_2^2 + h(x^{k+1})$$

$$408 \quad - G(x^*, x^k) \leq \epsilon_2^k - \frac{1}{2s} \|x^* - \bar{x}^{k+1}\|_2^2$$

$$409 \quad (4.13) \iff g(x^k) + \nabla g(x^k)^\top (x^{k+1} - x^k) + \epsilon_1^{k\top} (x^{k+1} - x^k)$$

$$410 \quad + \frac{1}{2s} \|x^{k+1} - x^k\|_2^2 + h(x^{k+1}) - G(x^*, x^k) \leq \epsilon_2^k - \frac{1}{2s} \|x^* - \bar{x}^{k+1}\|_2^2.$$

412 From (4.11) to (4.12), we simply used the definition of  $G(x, x^k)$  in (4.6) with  $x = x^{k+1}$   
 413 and we also used  $\nabla^{\epsilon_1^k} g(x^k) := \nabla g(x^k) + \epsilon_1^k$  in (4.13).

414 Applying (3.4) to (4.13) (with  $s \leq 1/L$ ) and using  $f := g + h$ , we obtain

$$415 \quad (4.14) \quad g(x^{k+1}) + h(x^{k+1}) - G(x^*, x^k) \leq \epsilon_2^k - \frac{1}{2s} \|x^* - \bar{x}^{k+1}\|_2^2$$

$$416 \quad + \epsilon_1^{k\top} (x^k - x^{k+1}),$$

$$417 \quad (4.15) \iff f(x^{k+1}) - G(x^*, x^k) \leq \epsilon_2^k - \frac{1}{2s} \|x^* - \bar{x}^{k+1}\|_2^2 + \epsilon_1^{k\top} (x^k - x^{k+1}).$$

419 We now expand  $G(x^*, x^k)$  in (4.15) as follows

$$420 \quad (4.16) \quad \begin{aligned} f(x^{k+1}) - g(x^k) - \nabla^{\epsilon_1^k} g(x^k)^\top (x^* - x^k) - \frac{1}{2s} \|x^* - x^k\|_2^2 - h(x^*) \\ \leq \epsilon_2^k - \frac{1}{2s} \|x^* - \bar{x}^{k+1}\|_2^2 + \epsilon_1^{k\top} (x^k - x^{k+1}). \end{aligned}$$

421 Rearranging and subtracting  $g(x^*)$  from both sides yields

$$422 \quad (4.17) \quad \begin{aligned} f(x^{k+1}) - h(x^*) - g(x^*) \leq -g(x^*) + \epsilon_2^k - \frac{1}{2s} \|x^* - \bar{x}^{k+1}\|_2^2 + g(x^k) \\ + \nabla^{\epsilon_1^k} g(x^k)^\top (x^* - x^k) + \frac{1}{2s} \|x^* - x^k\|_2^2 + \epsilon_1^{k\top} (x^k - x^{k+1}). \end{aligned}$$

423 Using the definitions  $f := g + h$  and  $\nabla^{\epsilon_1^k} g(x^k) = \nabla g(x^k) + \epsilon_1^k$  in (4.17), we obtain

$$424 \quad (4.18) \quad \begin{aligned} f(x^{k+1}) - f(x^*) \leq \epsilon_2^k - g(x^*) + g(x^k) + \nabla g(x^k)^\top (x^* - x^k) \\ - \frac{1}{2s} \|x^* - \bar{x}^{k+1}\|_2^2 + \frac{1}{2s} \|x^* - x^k\|_2^2 + \epsilon_1^{k\top} (x^* - x^k) + \epsilon_1^{k\top} (x^k - x^{k+1}) \\ \leq \epsilon_2^k - \frac{1}{2s} \|x^* - \bar{x}^{k+1}\|_2^2 + \frac{1}{2s} \|x^* - x^k\|_2^2 + \epsilon_1^{k\top} (x^* - x^{k+1}), \end{aligned}$$

425 where in the second inequality we used the fact that  $g$  is convex, i.e.,  $g(x^*) \geq g(x^k) +$   
 426  $\nabla g(x^k)^\top(x^* - x^k)$ . Summing both sides of (4.18) from 0 to  $k$ ,

$$\begin{aligned}
 (4.19) \quad & \sum_{i=0}^k [f(x^{i+1}) - f(x^*)] \leq \sum_{i=0}^k \epsilon_2^i + \sum_{i=0}^k \epsilon_1^{i\top} (x^* - x^{i+1}) \\
 & \quad + \frac{1}{2s} \sum_{i=0}^k \left[ \|x^* - x^i\|_2^2 - \|x^* - \bar{x}^{i+1}\|_2^2 \right], \\
 & = \sum_{i=0}^k \epsilon_2^i + \sum_{i=0}^k \epsilon_1^{i\top} (x^* - x^{i+1}) + \frac{1}{2s} \sum_{i=0}^k \left[ \|x^* - x^i\|_2^2 \right. \\
 & \quad \left. - (\|x^* - x^{i+1}\|_2^2 + \|x^{i+1} - \bar{x}^{i+1}\|_2^2 \right. \\
 427 & \quad \left. + 2(x^{i+1} - \bar{x}^{i+1})^\top(x^* - x^{i+1})) \right], \\
 & = \sum_{i=0}^k \epsilon_2^i + \sum_{i=0}^k \epsilon_1^{i\top} (x^* - x^{i+1}) + \frac{1}{2s} \sum_{i=0}^k \left[ \|x^* - x^i\|_2^2 \right. \\
 & \quad \left. - (\|x^* - x^{i+1}\|_2^2 + \|r^{i+1}\|_2^2 + 2(r^{i+1})^\top(x^* - x^{i+1})) \right], \\
 & = \sum_{i=0}^k \epsilon_2^i + \sum_{i=0}^k (\epsilon_1^i - \frac{1}{s} r^{i+1})^\top (x^* - x^{i+1}) + \frac{1}{2s} \left[ \|x^* - x^0\|_2^2 \right. \\
 & \quad \left. - \|x^* - x^{k+1}\|_2^2 \right] - \frac{1}{2s} \sum_{i=0}^k \|r^{i+1}\|_2^2,
 \end{aligned}$$

428 where in the second-to-last equality we used the definition of  $r^i$  in (3.11), and in the  
 429 last equality we noticed that the quadratic terms involving  $x^*$  formed a telescopic  
 430 sequence. Rearranging and moving negative terms to the left hand side results in

$$\begin{aligned}
 (4.20) \quad & \sum_{i=0}^k [f(x^{i+1}) - f(x^*)] + \frac{1}{2s} \sum_{i=0}^k \|r^{i+1}\|_2^2 + \frac{1}{2s} \|x^* - x^{k+1}\|_2^2 \leq \sum_{i=0}^k \epsilon_2^i \\
 431 & \quad + \sum_{i=0}^k (\epsilon_1^i - \frac{1}{s} r^{i+1})^\top (x^* - x^{i+1}) + \frac{1}{2s} \|x^* - x^0\|_2^2.
 \end{aligned}$$

432 Since  $f$  is a convex function, Jensen's inequality implies

$$433 \quad f\left(\frac{1}{k+1} \sum_{i=0}^k x^{i+1}\right) - f(x^*) \leq \frac{1}{k+1} \sum_{i=0}^k [f(x^{i+1}) - f(x^*)],$$

434 which, applied to (4.20) and together with the fact that the last two terms of the  
 435 left-hand side of (4.20) are nonnegative, yields

$$\begin{aligned}
 436 \quad & f\left(\frac{1}{k+1} \sum_{i=0}^k x^{i+1}\right) - f(x^*) + \frac{1}{2(k+1)s} \sum_{i=0}^k \|r^{i+1}\|_2^2 + \frac{1}{2(k+1)s} \|x^* - x^{k+1}\|_2^2 \leq \\
 437 & \quad \frac{1}{k+1} \left[ \sum_{i=0}^k \epsilon_2^i + \sum_{i=0}^k (\epsilon_1^i - \frac{1}{s} r^{i+1})^\top (x^* - x^{i+1}) + \frac{1}{2s} \|x^* - x^0\|_2^2 \right]. \\
 438 &
 \end{aligned}$$

439 Using Lemma 1 to bound the norm of the residual error  $r^k = x^k - \bar{x}^k$  resulting  
 440 from the proximal error  $\epsilon_2^k$ , Cauchy-Schwarz yields

$$\begin{aligned}
 441 \quad (4.22) \quad (\epsilon_1^i - \frac{1}{s}r^{i+1})^\top(x^* - x^{i+1}) &\leq \left( \|\epsilon_1^i\|_2 + \frac{1}{s} \|r^{i+1}\|_2 \right) \|x^* - x^{i+1}\|_2 \\
 &\leq \left( \|\epsilon_1^i\|_2 + \sqrt{\frac{2\epsilon_2^i}{s}} \right) \|x^* - x^{i+1}\|_2.
 \end{aligned}$$

442 Using (4.22) in (4.21) yields

$$\begin{aligned}
 443 \quad (4.23) \quad f\left(\frac{1}{k+1} \sum_{i=0}^k x^{i+1}\right) - f(x^*) &\leq \frac{1}{k+1} \sum_{i=0}^k \epsilon_2^i \\
 &\quad + \frac{1}{k+1} \sum_{i=0}^k \left( \|\epsilon_1^i\|_2 + \sqrt{\frac{2\epsilon_2^i}{s}} \right) \|x^* - x^{i+1}\|_2 \\
 &\quad + \frac{1}{2s(k+1)} \|x^* - x^0\|_2^2
 \end{aligned}$$

444 Applying Quasi-Féjer (Theorem 6 in the appendix) recursively gives

$$\begin{aligned}
 445 \quad (4.24) \quad f\left(\frac{1}{k+1} \sum_{i=0}^k x^{i+1}\right) - f(x^*) &\leq \frac{1}{k+1} \sum_{i=0}^k \epsilon_2^i + \frac{1}{2s(k+1)} \|x^* - x^0\|_2^2 \\
 &\quad + \frac{1}{k+1} \sum_{i=0}^k \left( \|\epsilon_1^i\|_2 + \sqrt{\frac{2\epsilon_2^i}{s}} \right) \|x^* - x^0\|_2 \\
 &\quad + \frac{1}{k+1} \sum_{i=0}^k \left( \|\epsilon_1^i\|_2 + \sqrt{\frac{2\epsilon_2^i}{s}} \right) \left( \sum_{j=1}^i E^j + iC_\rho \right),
 \end{aligned}$$

446 where  $E^j = \|r^j\|_2 + s_{j-1} \|\epsilon_1^{j-1}\|_2$  and  $C_\rho = 0$  if the optimum  $x^*$  is reached. This  
 447 completes the proof of Theorem 1.

448 **4.2. Proof of Theorem 2.** This result is about the basic version of approximate  
 449 PGD, but with random proximal computation error  $\epsilon_{2\Omega}$ , component-wise bounded  
 450 gradient error  $\epsilon_{1\Omega}$  and bounded residuals  $\|x_\Omega^k - x^*\|_2$ . As the algorithm generates a  
 451 sequence of random vectors  $\{x_\Omega^k\}$ , the residual vector sequence  $\{r_\Omega^k\}$  will also be a  
 452 random.

453 Let  $T_k$  denote the second error term in the bound of (3.14) [Theorem 1], i.e.,

$$454 \quad (4.25) \quad T_k = \begin{cases} 0 & , k = 0 \\ \sum_{i=1}^k (\epsilon_{1\Omega}^{i-1} - \frac{1}{s}r_\Omega^i)^\top(x^* - x_\Omega^i) & , k = 1, 2, \dots, \end{cases}$$

455 The first step is to show that  $\{T_k\}$  is a martingale. Recall that a sequence of random  
 456 variables  $T_0, T_1, \dots$  is a martingale with respect to the sequence  $X_0, X_1, \dots$  if, for all  
 457  $k \geq 0$ , the following conditions hold:

- 458 •  $T_k$  is a function of  $X_0, X_1, \dots, X_k$ ;
- 459 •  $\mathbb{E}[|T_k|] < \infty$ ;
- 460 •  $\mathbb{E}[T_{k+1}|X_0, X_1, \dots, X_k] = T_k$ .

461 A sequence of random variables  $T_0, T_1, \dots$  is called a martingale when it is a martin-  
 462 gale with respect to itself. That is,  $\mathbb{E}[|T_k|] < \infty$ , and  $\mathbb{E}[T_{k+1}|T_0, T_1, \dots, T_k] = T_k$ .

463 Let  $\nu_\Omega^k = \epsilon_{1\Omega}^{k-1} - \frac{1}{s}r_\Omega^k$  and recall the definition of  $r_\Omega^k$  in (3.11):

$$464 \quad (4.26) \quad r^k = x^k - \bar{x}^k.$$

465 Rewriting (4.25) in terms of  $\nu_\Omega^k$  yields

$$466 \quad (4.27) \quad T_k = T_{k-1} + \nu_\Omega^{k\top} (x^* - x_\Omega^k).$$

467 We now show that Assumptions 2 and 3 imply that  $\{T_k\}_{k \geq 0}$  is a martingale. Specif-  
468 ically, (3.10a) and (3.13a), we have

$$469 \quad \mathbb{E}[\nu_\Omega^k | \nu_\Omega^1 \dots \nu_\Omega^{k-1}] = \mathbb{E}[\nu_\Omega^k] = 0.$$

470 And from (3.10c) and (3.13b), we have

$$471 \quad \mathbb{E}[\nu_\Omega^{k\top} x_\Omega^k | \nu_\Omega^1 \dots \nu_\Omega^{k-1}, x_\Omega^1 \dots x_\Omega^{k-1}] = \mathbb{E}[\nu_\Omega^{k\top} x_\Omega^k] = 0.$$

472 Taking the expected value of both sides of (4.27) conditioned on  $\{T_i\}_{i=1}^{k-1}$  gives

$$\begin{aligned} 473 \quad \mathbb{E}[T_k | T_1 \dots T_{k-1}] &= \mathbb{E}[T_{k-1} + \nu_\Omega^{k\top} (x^* - x_\Omega^k) | T_1 \dots T_{k-1}] \\ 474 &= \mathbb{E}[T_{k-1} | T_1 \dots T_{k-1}] + \mathbb{E}[\nu_\Omega^{k\top} (x^* - x_\Omega^k) | T_1 \dots T_{k-1}] \\ 475 &= T_{k-1} + \mathbb{E}[\nu_\Omega^{k\top} (x^* - x_\Omega^k) | T_1 \dots T_{k-1}] \\ 476 &= T_{k-1} + \mathbb{E}[\nu_\Omega^{k\top} x^* | T_1 \dots T_{k-1}] - \mathbb{E}[\nu_\Omega^{k\top} x_\Omega^k | T_1 \dots T_{k-1}] \\ 477 &= T_{k-1} + \mathbb{E}[\nu_\Omega^k | T_1 \dots T_{k-1}]^\top x^* - \mathbb{E}[\nu_\Omega^{k\top} x_\Omega^k | T_1 \dots T_{k-1}] \\ 478 \quad (4.28) &= T_{k-1} + \mathbb{E}[\nu_\Omega^k | \nu_\Omega^1 \dots \nu_\Omega^{k-1}, x_\Omega^1 \dots x_\Omega^{k-1}]^\top x^* \\ 479 &\quad - \mathbb{E}[\nu_\Omega^{k\top} x_\Omega^k | \nu_\Omega^1 \dots \nu_\Omega^{k-1}, x_\Omega^1 \dots x_\Omega^{k-1}] \end{aligned}$$

$$480 \quad (4.29) \quad = T_{k-1} + \mathbb{E}\left[\left(\epsilon_{1\Omega}^{k-1} - \frac{1}{s}r_\Omega^k\right)^\top x^* - \mathbb{E}\left[\left(\epsilon_{1\Omega}^{k-1} - \frac{1}{s}r_\Omega^k\right)^\top x_\Omega^k\right]\right]$$

$$481 \quad (4.30) \quad = T_{k-1} + \mathbb{E}\left[\left(\epsilon_{1\Omega}^{k-1} - \frac{1}{s}r_\Omega^k\right)^\top x^* - \mathbb{E}\left[\mathbb{E}\left[\left(\epsilon_{1\Omega}^{k-1} - \frac{1}{s}r_\Omega^k\right)^\top x_\Omega^k\right] \middle| x_\Omega^k\right]\right]$$

$$482 \quad (4.31) \quad = T_{k-1} - \mathbb{E}\left[\left(\epsilon_{1\Omega}^{k-1} - \frac{1}{s}r_\Omega^k\right)^\top x_\Omega^k\right]$$

$$483 \quad (4.32) \quad = T_{k-1}.$$

485 From (4.28) to (4.29), we used the error mean independence assumption, i.e.,  
486  $\mathbb{E}[\nu_\Omega^k | \nu_\Omega^1 \dots \nu_\Omega^{k-1}] = \mathbb{E}[\nu_\Omega^k]$  as well as the data mean independence assumption (or  
487 the less restrictive statistical orthogonality in high dimensional problems), i.e.,

488  $\mathbb{E}[\nu_\Omega^{k\top} x_\Omega^k | \nu_\Omega^1 \dots \nu_\Omega^{k-1}, x_\Omega^1 \dots x_\Omega^{k-1}] = \mathbb{E}[\nu_\Omega^{k\top} x_\Omega^k]$ . From (4.31) to (4.32), we used the  
489 zero mean error assumption, i.e.,  $\mathbb{E}[\nu_\Omega^k] = 0$ . Therefore,  $T_1, T_2, \dots, T_k$  is a martingale.

490 In what follows, we establish upper bounds on the absolute value of the martingale  
491  $\{T_k\}$ . To do that, we use the Azuma-Hoeffding inequality in [27, p. 36], noticing that  
492  $|T_k - T_{k-1}| = |\nu_\Omega^{k\top} (x^* - x_\Omega^k)| \leq \left(\sqrt{n}\delta M_{\nabla g} + \sqrt{2\epsilon_2^k/s}\right) \|x_\Omega^* - x_\Omega^k\|_2$ , where we have  
493 used Cauchy-Schwarz, etc. Corollary [27, Corollary 2.20] then yields

$$494 \quad (4.33) \quad \Pr\left(|T_k - T_0| \geq \gamma \sqrt{\sum_{i=1}^k \left(\sqrt{n}M_{\nabla g}|\delta| + \sqrt{\frac{2\epsilon_2^i}{s}}\right)^2 \|x_\Omega^* - x_\Omega^i\|_2^2}\right) \leq 2 \exp\left(-\frac{\gamma^2}{2}\right).$$

495 Since  $\epsilon_2^k \leq \epsilon_0$ , then the following also holds

$$496 \quad (4.34) \quad \Pr\left(|T_k - T_0| \geq \gamma(\sqrt{n}M_{\nabla g}|\delta| + \sqrt{\frac{2\epsilon_0}{s}})\sqrt{\sum_{i=1}^k \|x_\Omega^* - x_\Omega^i\|_2^2}\right) \leq 2\exp(-\frac{\gamma^2}{2}).$$

497 And since  $T_0 = 0$  we obtain

$$498 \quad (4.35) \quad \Pr\left(|T_k| \geq \gamma(\sqrt{n}M_{\nabla g}|\delta| + \sqrt{\frac{2\epsilon_0}{s}})\sqrt{\sum_{i=1}^k \|x_\Omega^* - x_\Omega^i\|_2^2}\right) \leq 2\exp(-\frac{\gamma^2}{2}).$$

499 Or, equivalently, that

$$500 \quad (4.36) \quad |T_k| \leq \gamma(\sqrt{n}M_{\nabla g}|\delta| + \sqrt{\frac{2\epsilon_0}{s}})\sqrt{\sum_{i=1}^k \|x_\Omega^* - x_\Omega^i\|_2^2}$$

501 holds for all  $k \geq 1$  with probability at least  $1 - 2\exp(-\frac{\gamma^2}{2})$ . Expanding  $T_k$  we obtain

$$502 \quad (4.37) \quad \left|\sum_{i=1}^k (\epsilon_{1\Omega}^{i-1} - \frac{1}{s}r_\Omega^i)^\top (x_\Omega^* - x_\Omega^i)\right| \leq \gamma\left(\sqrt{n}M_{\nabla g}|\delta| + \sqrt{\frac{2\epsilon_0}{s}}\right)\sqrt{\sum_{i=1}^k \|x_\Omega^* - x_\Omega^i\|_2^2}.$$

503 By assumption, we have that  $\|x_\Omega^* - x_\Omega^i\|_2^2 \leq D_x \|x_\Omega^* - x_\Omega^0\|_2^2$  holds with probability  $p$ ,  
504 for each  $i$ . Then,

$$505 \quad (4.38) \quad \left|\sum_{i=1}^k (\epsilon_{1\Omega}^{i-1} - \frac{1}{s}r_\Omega^i)^\top (x_\Omega^* - x_\Omega^i)\right| \leq \gamma\left(M_{\nabla g}\sqrt{nk}|\delta| + \sqrt{\frac{2k\epsilon_0}{s}}\right)D_x \|x_\Omega^* - x_\Omega^0\|_2$$

506 holds with probability  $p^k(1 - 2\exp(-\frac{\gamma^2}{2}))$ . Substituting (4.38) into (3.14) completes  
507 the proof of Theorem 2.

508 **4.3. Proof of Theorem 3.** Here  $\epsilon_{2\Omega}$  is bounded almost surely and has station-  
509 ary mean. Specifically, we have  $0 \leq \epsilon_{2\Omega}^k \leq \epsilon_0$ , with probability 1. By Hoeffding's  
510 inequality ([27, Proposition 2.5]), we can write,

$$511 \quad (4.39) \quad \Pr\left(\left|\sum_{i=1}^k \epsilon_{2\Omega}^i - \mathbb{E}\left(\sum_{i=1}^k \epsilon_{2\Omega}^i\right)\right| \geq t\right) \leq 2\exp\left(\frac{-2t^2}{k\epsilon_0^2}\right), \quad \text{for all } t > 0.$$

512 Defining the constant mean  $\mathbb{E}[\epsilon_{2\Omega}^k] = \mathbb{E}[\epsilon_{2\Omega}]$  and substituting in (4.39) yields

$$513 \quad (4.40) \quad \Pr\left(\left|\sum_{i=1}^k \epsilon_{2\Omega}^i - k\mathbb{E}[\epsilon_{2\Omega}]\right| \geq t\right) \leq 2\exp\left(\frac{-2t^2}{k\epsilon_0^2}\right), \quad \text{for all } t > 0.$$

514 By choosing  $t = \frac{\gamma\sqrt{k\epsilon_0}}{2}$ , for some  $\gamma > 0$ , we obtain

$$515 \quad (4.41) \quad \Pr\left(\left|\sum_{i=1}^k \epsilon_{2\Omega}^i - k\mathbb{E}[\epsilon_{2\Omega}]\right| \geq \frac{\gamma\sqrt{k\epsilon_0}}{2}\right) \leq 2\exp\left(\frac{-\gamma^2}{2}\right) \quad \text{for all } \gamma > 0.$$



516 Equivalently,

$$517 \quad (4.42) \quad \sum_{i=1}^k \epsilon_{2\Omega}^i \leq k\mathbb{E}[\epsilon_{2\Omega}] + \frac{\gamma\sqrt{k}\epsilon_0}{2}$$

518 holds with probability at least  $1 - 2\exp(-\frac{\gamma^2}{2})$ . Using the last inequality (4.42) in  
519 (3.16) and applying the probability union bound completes the proof of Theorem 3.

520 **4.4. Proof of Theorem 4.** Following the same line of proof of Section 4.1 but  
521 with  $y_k = (1 + \beta_k)x^k - \beta_k x^{k-1}$ , where  $\{\beta_k\} \in [0, 1]$  is the momentum sequence, and  
522 using the approximate accelerated PG iteration scheme 3.6, we obtain

$$523 \quad (4.43) \quad \begin{aligned} f(x^{k+1}) - f(x) &\leq \epsilon_2^k + \epsilon_1^{k\top} (x - x^{k+1}) - \frac{1}{2s} \|x - x^{k+1}\|_2^2 \\ &\quad - \frac{1}{2s} (r^{k+1})^\top (x - x^{k+1}) + \frac{1}{2s} \|x - y^k\|_2^2. \end{aligned}$$

524 Let us now substitute  $y^k$  and  $x$  by,

$$525 \quad (4.44) \quad y^k = x^k + \beta_k(x^k - x^{k-1})$$

$$526 \quad (4.45) \quad x = \alpha_k^{-1}x^* + (1 - \alpha_k^{-1})x^k,$$

528 where (4.44) follows from the definition of the acceleration scheme (3.6), and (4.45)  
529 is a choice that we make to simplify the analysis.<sup>4</sup>  $\{\alpha_k\}_{k \geq 1}$  is a given parameter  
530 sequence that satisfies  $\alpha_0 = 1$ ,  $\alpha_k \geq 1$  and  $\beta_k = \frac{\alpha_{k-1}-1}{\alpha_k}$ . (4.43) can now be expanded  
531 as

$$532 \quad (4.46) \quad \begin{aligned} f(x^{k+1}) - f(\alpha_k^{-1}x^* + (1 - \alpha_k^{-1})x^k) &\leq \epsilon_2^k + \epsilon_1^{k\top} (\alpha_k^{-1}x^* + (1 - \alpha_k^{-1})x^k - x^{k+1}) \\ &\quad - \frac{1}{2s} \|\alpha_k^{-1}x^* + (1 - \alpha_k^{-1})x^k - x^{k+1}\|_2^2 \\ &\quad + \frac{1}{2s} \|\alpha_k^{-1}x^* + (1 - \alpha_k^{-1})x^k - y^k\|_2^2 \\ &\quad - \frac{1}{2s} (r^{k+1})^\top (\alpha_k^{-1}x^* + (1 - \alpha_k^{-1})x^k - x^{k+1}). \end{aligned}$$

533 Since  $\alpha_k^{-1} \in ]0, 1]$ , and from the convexity of  $f$ , we have

$$534 \quad (4.47) \quad \begin{aligned} f(x^{k+1}) - f(\alpha_k^{-1}x^* + (1 - \alpha_k^{-1})x^k) &\geq f(x^{k+1}) + (1 - \alpha_k^{-1})f(x^*) \\ &\quad - (1 - \alpha_k^{-1})f(x^k) - f(x^*) \\ &= f(x^{k+1}) - f(x^*) - (1 - \alpha_k^{-1})(f(x^k) - f(x^*)). \end{aligned}$$

535 Let us now define the new sequences  $\{v^k\}$  and  $\{u^k\}$  by

$$536 \quad (4.48) \quad u^k := x^* + (\alpha_k - 1)x^k - \alpha_k y^k = x^* - (x^k + (\alpha_{k-1} - 1)(x^k - x^{k-1}))$$

$$537 \quad (4.49) \quad v^k = f(x^k) - f(x^*).$$

<sup>4</sup>Note that  $y^k \rightarrow x^k$  as  $x^k \rightarrow x^*$ .

539 From these we can obtain

$$540 \quad (4.50) \quad u^{k+1} := x^* + (\alpha_k - 1)x^k - \alpha_k x^{k+1} = x^* - (x^{k+1} + (\alpha_k - 1)(x^{k+1} - x^k)),$$

542 by using  $\beta_k = (\alpha_{k-1} - 1)/\alpha_k$  and  $y^k = (1 + \beta_k)x^k - \beta_k x^{k-1}$ .

543 Rewriting (4.46) in terms of the newly defined sequences,  $\{u^k\}$  and  $\{v^k\}$ , and

544 using (4.47) with  $c_k := 1 - \alpha^{-1}$ , as well as (4.48) and (4.50) we obtain

$$545 \quad (4.51) \quad \begin{aligned} v^{k+1} - c_k v^k &\leq \epsilon_2^k + \frac{1}{\alpha_k} \epsilon_1^{k\top} u^{k+1} - \frac{1}{2s\alpha_k^2} \|u^{k+1}\|_2^2 + \frac{1}{2s\alpha_k^2} \|u^k\|_2^2 \\ &\quad - \frac{1}{2s} \|r^{k+1}\|_2^2 - \frac{1}{2s\alpha_k} (r^{k+1})^\top u^{k+1}. \end{aligned}$$

546 Rearranging (4.51) we obtain

$$547 \quad (4.52) \quad \begin{aligned} v^{k+1} + \frac{1}{2s} \|r^{k+1}\|_2^2 + \frac{1}{2s\alpha_k^2} \|u^{k+1}\|_2^2 &\leq \epsilon_2^k + \frac{1}{\alpha_k} \epsilon_1^{k\top} u^{k+1} + c_k v^k \\ &\quad + \frac{1}{2s\alpha_k^2} \|u^k\|_2^2 - \frac{1}{2s\alpha_k} (r^{k+1})^\top u^{k+1}. \end{aligned}$$

548 Multiplying both sides by  $\alpha_k^2$ ,

$$549 \quad (4.53) \quad \begin{aligned} \alpha_k^2 v^{k+1} + \frac{\alpha_k^2}{2s} \|r^{k+1}\|_2^2 + \frac{1}{2s} \|u^{k+1}\|_2^2 &\leq \alpha_k^2 \epsilon_2^k + \alpha_k \epsilon_1^{k\top} u^{k+1} + \alpha_k^2 c_k v^k \\ &\quad + \frac{1}{2s} \|u^k\|_2^2 - \frac{\alpha_k}{2s} (r^{k+1})^\top u^{k+1}. \end{aligned}$$

550 Applying (4.53) recursively, and substituting  $\alpha_k^2 c_k = \alpha_k^2 - \alpha_k = \alpha_{k-1}$  yields

$$551 \quad (4.54) \quad \begin{aligned} \alpha_k^2 v^{k+1} + \frac{\alpha_k^2}{2s} \|r^{k+1}\|_2^2 + \frac{1}{2s} \|u^{k+1}\|_2^2 &\leq \alpha_k^2 \epsilon_2^k + \alpha_k \epsilon_1^{k\top} u^{k+1} + \alpha_{k-1} v^k \\ &\quad + \frac{1}{2s} \|u^k\|_2^2 - \frac{\alpha_k}{2s} (r^{k+1})^\top u^{k+1}, \end{aligned}$$

552

...

$$554 \quad (4.55) \quad \begin{aligned} \alpha_1^2 v^2 + \frac{\alpha_1^2}{2s} \|r^2\|_2^2 + \frac{1}{2s} \|u^2\|_2^2 &\leq \alpha_1^2 \epsilon_2^2 + \alpha_1 \epsilon_1^{2\top} u^2 + \alpha_0 v^1 \\ &\quad + \frac{1}{2s} \|u^1\|_2^2 - \frac{\alpha_1}{2s} (r^2)^\top u^2. \end{aligned}$$

555

556 Adding both sides of all inequalities,

$$558 \quad (4.56) \quad \begin{aligned} \alpha_k^2 v^{k+1} + \sum_{i=0}^k \frac{\alpha_i^2}{2s} \|r^{i+1}\|_2^2 + \frac{1}{2s} \|u^{k+1}\|_2^2 &+ \sum_{i=0}^k (\alpha_{i-1}^2 - \alpha_{i-1}) v^i \\ &\leq \sum_{i=0}^k \alpha_i^2 \epsilon_2^i + \frac{1}{2s} \|u^1\|_2^2 + \sum_{i=0}^k \alpha_i \epsilon_1^{i\top} u^{i+1} + \alpha_0 v^1 - \sum_{i=0}^k \frac{\alpha_i}{2s} (r^{i+1})^\top u^{i+1}. \end{aligned}$$

559 Substituting  $\alpha_{i-1}^2 - \alpha_{i-1} = \alpha_{i-2}^2$  and  $\alpha_0 = 1$  gives,

$$560 \quad \begin{aligned} \alpha_k^2 v^{k+1} + \sum_{i=0}^k \frac{\alpha_i^2}{2s} \|r^{i+1}\|_2^2 + \frac{1}{2s} \|u^{k+1}\|_2^2 &+ \sum_{i=0}^k \alpha_{i-2} v^i \\ &\leq \sum_{i=0}^k \alpha_i^2 \epsilon_2^i + \sum_{i=0}^k \alpha_i \epsilon_1^{i\top} u^{i+1} + v^1 + \frac{1}{2s} \|u^1\|_2^2 - \sum_{i=0}^k \frac{\alpha_i}{2s} (r^{i+1})^\top u^{i+1}. \end{aligned}$$

561 For a positive sequence  $\{\alpha_k\}_{k \geq 0}$  and because  $x^*$  is a (global) minimizer,  $\sum \alpha_{i-2} v^i \geq 0$   
 562 is always satisfied; hence the following holds

$$\begin{aligned}
 \alpha_k^2 v^{k+1} &\leq \alpha_k^2 v^{k+1} + \sum_{i=0}^k \frac{\alpha_i^2}{2s} \|r^{i+1}\|_2^2 + \frac{1}{2s} \|u^{k+1}\|_2^2 + \sum_{i=0}^k \alpha_{i-2} v^i \\
 (4.57) \quad &\leq \sum_{i=0}^k \alpha_i^2 \epsilon_2^i + \sum_{i=0}^k \alpha_i \left( \epsilon_1^i - \frac{1}{s} r^{i+1} \right)^\top u^{i+1} + v^1 + \frac{1}{2s} \|u^1\|_2^2.
 \end{aligned}$$

564 From (4.43) with  $k = 0$  and  $x = x^*$ , we have

$$\begin{aligned}
 (4.58) \quad v^1 = f(x^1) - f(x^*) &\leq \epsilon_2^0 + \left( \epsilon_1^0 - \frac{1}{2s} r^1 \right)^\top (x^* - x^1) - \frac{1}{2s} \|x^* - x^1\|_2^2 \\
 &\quad + \frac{1}{2s} \|x^* - x^0\|_2^2,
 \end{aligned}$$

566 since  $y^0 = x^0$ . From the definition of  $\{u^k\}$  in (4.50) we have

$$\begin{aligned}
 (4.59) \quad \frac{1}{2s} \|u^1\|_2^2 &= \frac{1}{2s} \|x^* + (\alpha_0 - 1)x^0 - \alpha_0 x^1\|_2^2, \\
 &= \frac{1}{2s} \|x^* - x^1\|_2^2,
 \end{aligned}$$

568 where we have used the initialization  $\alpha_0 = 1$ . Substituting for  $v^{k+1}$  and combining  
 569 (4.58) and (4.59) with (4.57) yields

$$(4.60) \quad \alpha_k^2 (f(x^{k+1}) - f(x^*)) \leq \sum_{i=0}^k \alpha_i^2 \epsilon_2^i + \sum_{i=0}^k \alpha_i \left( \epsilon_1^i - \frac{1}{s} r^{i+1} \right)^\top u^{i+1} + \frac{1}{2s} \|x^* - x^0\|_2^2.$$

571 Dividing both sides by  $\alpha_k^2$  and applying Cauchy-Schwarz inequality yields

$$\begin{aligned}
 (4.61) \quad f(x^{k+1}) - f(x^*) &\leq \frac{1}{\alpha_k^2} \left[ \sum_{i=0}^k \alpha_i^2 \epsilon_2^i + \left[ \sum_{i=0}^k \alpha_i \left( \|\epsilon_1^i\|_2 + \frac{1}{s} \|r^{i+1}\|_2 \right) \right] \|u^{i+1}\|_2 \right. \\
 &\quad \left. + \frac{1}{s} \|x^* - x^0\|_2^2 \right].
 \end{aligned}$$

575 We have by definition 4.48 and 4.50

$$(4.62) \quad u^k = x^* + (\alpha_k - 1)x^k - \alpha_k y^k = x^* - (x^k + (\alpha_{k-1} - 1)(x^k - x^{k-1})),$$

$$(4.63) \quad u^{k+1} = x^* + (\alpha_k - 1)x^k - \alpha_k x^{k+1} = x^* - (x^{k+1} + (\alpha_k - 1)(x^{k+1} - x^k)).$$

579 By triangle inequality of the vector norm, we have

$$\begin{aligned}
 (580) \quad \|u^k\|_2 &\leq \|(\alpha_k - 1)(x^k - x^*)\|_2 + \alpha_k \|y^k - x^*\|_2, \\
 (581) \quad \|u^{k+1}\|_2 &\leq |\alpha_k - 1| \|x^k - x^*\|_2 + \alpha_k \|x^{k+1} - x^*\|_2
 \end{aligned}$$

583 By the nonexpansivity of the displacement operator, i.e.,  $\mathbf{I} - s\nabla g$ , where  $\mathbf{I}$  is the  
 584 identity operator, we obtain

$$\begin{aligned}
 (4.64) \quad \|u^{k+1}\|_2 - \|u^k\|_2 &\leq \alpha_k \left| \|x^{k+1} - x^*\|_2 - \|y^k - x^*\|_2 \right|, \\
 &\leq \alpha_k \left| \|r^{k+1}\|_2 + s_k \|\epsilon_1^k\|_2 + C_{\rho, s_{k_0}} \right|, \quad \forall s_k \leq \frac{1}{L},
 \end{aligned}$$

588 where we have used inequality (A.18). Rearranging and taking into account that all  
589 the terms inside the absolute value are nonnegative, we obtain

$$590 \quad (4.65) \quad \|u^{k+1}\|_2 \leq \|u^k\|_2 + \alpha_k \left( \|r^{k+1}\|_2 + s_k \|\epsilon_1^k\|_2 + C_{\rho, s_{k_0}} \right), \quad \forall s_k \leq \frac{1}{L}.$$

592 Using the bound  $\|r^{i+1}\|_2 \leq \sqrt{2s\epsilon_2^i}$  from Lemma 1, by induction and backward sub-  
593 stitution

$$594 \quad (4.66) \quad \|u^{k+1}\|_2 \leq \|u^0\|_2 + \sum_{j=1}^k \alpha_j \left( \sqrt{2s\epsilon_2^j} + s_j \|\epsilon_1^j\|_2 + C_{\rho, s_{k_0}} \right), \quad \forall s_j \leq \frac{1}{L}.$$

596 where  $\|u^0\|_2 = \|x^0 - x^*\|_2$ . By multiplying we obtain the bound of Theorem 4.

597 **4.5. Proof of Theorem 5.** This result is about the accelerated version of ap-  
598 proximate PGD, but with random proximal computation error  $\epsilon_{2\Omega}$ , component-wise  
599 bounded gradient error  $\epsilon_{1\Omega}$  and bounded residuals  $\|x_\Omega^k - x^*\|_2$ . As the algorithm gen-  
600 erates a sequence of random vectors  $\{x_\Omega^k\}$ , the residual vector sequence  $\{r_\Omega^k\}$  will also  
601 be a random. Let  $\nu_\Omega = \epsilon_1^{i-1} - \frac{1}{s}r^i$  and let  $\{T_k\}$  denote the second error term in (3.14)  
602 [Theorem 4], i.e.,

$$603 \quad (4.67) \quad T_k = \begin{cases} 0, & k = 0 \\ \sum_{i=1}^k \alpha_i \nu_\Omega^i \top u_\Omega^i, & k = 1, 2, \dots, \end{cases}$$

604 where

$$605 \quad (4.68) \quad u_\Omega^i = x^* - x_\Omega^i + (1 - \alpha_{i-1})(x_\Omega^i - x_\Omega^{i-1}).$$

606 The first step is to show that  $\{T_k\}$  is a martingale. Recall that a sequence of random  
607 variables  $T_0, T_1, \dots$  is a martingale with respect to the sequence  $X_0, X_1, \dots$  if, for all  
608  $k \geq 0$ , the following conditions hold:

- 609 •  $T_k$  is a function of  $X_0, X_1, \dots, X_k$ ;
- 610 •  $\mathbb{E}[|T_k|] < \infty$ ;
- 611 •  $\mathbb{E}[T_{k+1} | X_0, X_1, \dots, X_k] = T_k$ .

612 A sequence of random variables  $T_0, T_1, \dots$  is called a martingale when it is a martin-  
613 gale with respect to itself. That is,  $\mathbb{E}[|T_k|] < \infty$ , and  $\mathbb{E}[T_{k+1} | T_0, T_1, \dots, T_k] = T_k$ . We  
614 now show that Assumptions 2 and 3 imply that  $\{T_k\}_{k \geq 0}$  is a martingale. Specifically,  
615 (3.10a) and (3.13a), we have

$$616 \quad \mathbb{E}[\nu_\Omega^k | \nu_\Omega^1 \dots \nu_\Omega^{k-1}] = \mathbb{E}[\nu_\Omega^k] = 0.$$

617 And from (3.10c) and (3.13b), we have

$$618 \quad \mathbb{E}[\nu_\Omega^k \top x_\Omega^k | \nu_\Omega^1 \dots \nu_\Omega^{k-1}, x_\Omega^1 \dots x_\Omega^{k-1}] = \mathbb{E}[\nu_\Omega^k \top x_\Omega^k] = 0.$$

619 We have from (4.67),

$$620 \quad (4.69) \quad T_k = T_{k-1} + \alpha_k \nu_\Omega^k \top u_\Omega^k.$$

621 Substituting for  $u_\Omega^k$  using (4.68) gives,

$$622 \quad (4.70) \quad T_k = T_{k-1} + \alpha_k \alpha_{k-1} \nu_\Omega^k \top (x^* - x_\Omega^k) + \alpha_k (1 - \alpha_{k-1}) \nu_\Omega^k \top (x^* - x_\Omega^{k-1}).$$

623 Taking the conditional expectation from both sides and proceeding as in Section 4.2,  
 624 we obtain  $\mathbb{E}[T_k | T_1 \dots T_{k-1}] = T_{k-1}$ , i.e.,  $T_1, T_2, \dots, T_k$  is a martingale.

625 In what follows, we establish upper bounds on the absolute value of the martingale  
 626  $\{T_k\}$ . By noticing that  $|T_k - T_{k-1}| = |\nu_\Omega^{k \top} u_\Omega^k| \leq \alpha_k (\sqrt{n} \delta M_{\nabla g} + \sqrt{2\epsilon_{2\Omega}^k/s}) \|u_\Omega^k\|_2$ ,  
 627 where we have used Cauchy-Schwarz, etc. [27, Corollary 2.20] then yields

$$\begin{aligned}
 628 \quad (4.71) \quad |T_k| &\leq \gamma |\delta| M_{\nabla g} \sqrt{n \sum_{i=1}^k i^2 \|u_\Omega^i\|_2^2} + \gamma \sqrt{2s} \sqrt{\sum_{i=1}^k i^2 \|u_\Omega^i\|_2^2 \epsilon_{2\Omega}^i} \\
 629 &\leq \gamma |\delta| M_{\nabla g} \sqrt{n} \sum_{i=1}^k i \|u_\Omega^i\|_2 + \gamma \sqrt{2s} \sum_{i=1}^k i \|u_\Omega^i\|_2 \sqrt{\epsilon_{2\Omega}^i} \\
 630
 \end{aligned}$$

631 where  $M_{\nabla g} = \sup_{i \in \mathbb{N}_+} \left\{ \|\nabla g(x^i)\|_\infty \right\}$  is the upper bound on the elements of the gradient.

632 Let  $\{S_k\}$  denote the first error term in (4) [Theorem 4] i.e.,

$$633 \quad (4.72) \quad S_k = \sum_{i=1}^k \alpha_i^2 \epsilon_{2\Omega}^i.$$

634 If  $0 \leq \epsilon_{2\Omega}^k \leq \epsilon_0$  and  $\alpha_k \leq k$ , then applying [27, Proposition 2.5] to  $S_k = \sum_{i=1}^k \alpha_i^2 \epsilon_{2\Omega}^i$   
 635 with  $0 \leq \epsilon_{2\Omega}^k \leq \epsilon_0$  and  $\alpha_k \leq k$  yields

$$636 \quad (4.73) \quad S_k \leq \mathbb{E} \left[ \sum_{i=1}^k \alpha_i^2 \epsilon_{2\Omega}^i \right] + \frac{\gamma}{2} \sqrt{\sum_{i=1}^k i^4 (\epsilon_{2\Omega}^i)^2} \leq \mathbb{E} \left[ \sum_{i=1}^k \alpha_i^2 \epsilon_{2\Omega}^i \right] + \frac{\gamma}{2} \sum_{i=1}^k i^2 \epsilon_{2\Omega}^i,$$

637 with probability at least  $1 - 2 \exp(-\frac{\gamma^2}{2})$ . Applying the probability union bound and  
 638 assuming that  $\|u_\Omega^i\|_2^2 \leq D_u^2 \|x^0 - x^*\|_2^2$  holds with probability  $p$  completes the proof  
 639 of Theorem 5.

640 **5. Experimental Results.** We now experimentally assess the proposed bounds  
 641 on an  $\ell_1$ -regularized model predictive control (MPC) problem. We consider a discrete  
 642 linear time invariant (LTI) state space model of a spacecraft [13]. The approximation  
 643 errors are simulated error sequences generated from a truncated Gaussian distribution.

644 **5.1. Model Predictive Control (MPC).** The  $\ell_1$ -regularized MPC can be  
 645 formulated as

$$646 \quad (5.1) \quad \underset{x \in \mathbb{R}^n}{\text{minimize}} \quad f(x) := g(x) + h(x),$$

647 where  $g : \mathbb{R}^n \rightarrow \mathbb{R}$  is the following real-valued, convex and differentiable function,

$$649 \quad g(x) := \left\| (\Phi^\top Q \Phi + R)^{\frac{1}{2}} x - (\Phi^\top Q \Phi + R)^{-\frac{1}{2}} \Phi^\top Q (R_s - \Psi x(k)) \right\|_2^2,$$

650 and  $h : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\}$  is the nondifferentiable convex  $\ell_1$ -norm

$$651 \quad h(x) := \lambda \|x\|_1,$$

652 with  $x \in \mathbb{R}^{p \times N_c \times 1}$  being the vectorized differential control  $\Delta u = u^k - u^{k-1} \in \mathbb{R}^{p \times N_c}$ ,  
 653 where  $p$  is the input dimension of the system and  $N_c$  is the control horizon. The

656 regularization parameter  $\lambda \in \mathbb{R}^+$  is a positive scalar.  $Q \in \mathbb{R}^{N_p \cdot m \times m \cdot N_p}$  and  $R \in$   
 657  $\mathbb{R}^{p \cdot N_c \times p \cdot N_c}$  are positive semi-definite design matrices where  $m$  is the output dimension  
 658 and  $N_p$  is the prediction horizon.  $R_s \in \mathbb{R}^{m \cdot N_p \times 1}$  is the vectorization of the matrix that  
 659 is constructed by  $N_p$  times stacking of the set-point vector  $r(k)$ .  $\Phi \in \mathbb{R}^{m \cdot N_p \times p \cdot N_c}$  and  
 660  $\Psi \in \mathbb{R}^{m \cdot N_p \times n}$  are augmented matrices which can be obtained from the spacecraft LTI  
 661 discrete state-space model  $(A, B, C)$  of [13] using a standard formula [28, Eq. 1.12].  
 662 For simulation, we select the problem's matrices as follows,

$$663 \quad Q = \text{diag}(500.0, 500.0, 500.0, 10^{-7}, 1.0, 1.0, 1.0, 500.0, 500.0, 500.0, 10^{-7}, 1.0, 1.0, 1.0);$$

$$664 \quad R = \text{diag}(200.0, 200.0, 200.0, 1.0, 200.0, 200.0, 200.0, 1.0),$$

666 and set the regularization parameter  $\lambda = 2.5021$ . The control and prediction horizons  
 667 are set to  $N_c = N_p = 5$ . The quadratic term of the  $\ell_1$ -regularized MPC problem,  
 668  $g(x)$ , has a gradient's Lipschitz constant of  $L = 11539$ , and therefore, a stepsize of  
 669  $s = \frac{1}{L}$  is adopted.

670 For the simulated errors, we use  $\epsilon_{1\Omega}^k = \nabla g(x^k) \odot \text{trand}(-\delta, \delta)$  where  $\text{trand}(a, b)$  is  
 671 the doubly truncated normal distribution [8] with lower and upper truncation points  
 672  $a$  and  $b$ , respectively.  $\delta$  is the gradient element-wise precision, which is a scalar upper  
 673 bound on the gradient error.  $\epsilon_2^k = \text{trand}(0, \epsilon_0)$  where  $\epsilon_0$  is a scalar upper bound on  
 674 the proximal computation error. The output of the distribution function  $\text{trand}(l, u)$   
 675 is a vector randomly generated from the standard multivariate normal distribution  
 676 truncated over the region  $[l, u]$ .

677 **5.2. Results.** The deterministic and probabilistic bounds are plotted and super-  
 678 imposed with the bound (2.1) of [25] in Figure 1 and Figure 2. The latter is denoted by  
 679 `Schmidt_1` (`Schmidt_2` in the accelerated case) and the proposed bounds are denoted  
 680 by `Thrm_1` and `Thrm_2` (`Thrm_4` and `Thrm_5` in the accelerated case), respectively.

681 Notice that we expect the effect of  $\epsilon_1^k$  to be negligible near the optimum since,  
 682 according to model (3.8),  $\epsilon_1^k$  is proportional to the magnitude of the gradient. How-  
 683 ever, depending on the choice of the upper bound of  $\epsilon_2^k$  in the proximal operation step  
 684 (3.7), the effect of the error  $\epsilon_2$  can still be significant and sometimes permanent even  
 685 near the optimum as we will see next.

686 In the presence of small gradient and proximal computation errors, the bounds  
 687 in Theorem 1, Theorem 2 practically coincide with (2.1). Therefore, in order to  
 688 emphasize the sharpness of the proposed bounds, we run the simulation with  $|\epsilon_1^k| \leq$   
 689  $2.2 \times 10$ ;  $\epsilon_2^k \leq 10$  for the nonaccelerated case (Figure 1), and with  $|\epsilon_1^k| \leq 2.2 \times 10^{-4}$ ;  $\epsilon_2^k \leq$   
 690  $10^{-4}$  for the accelerated case (Figure 2).

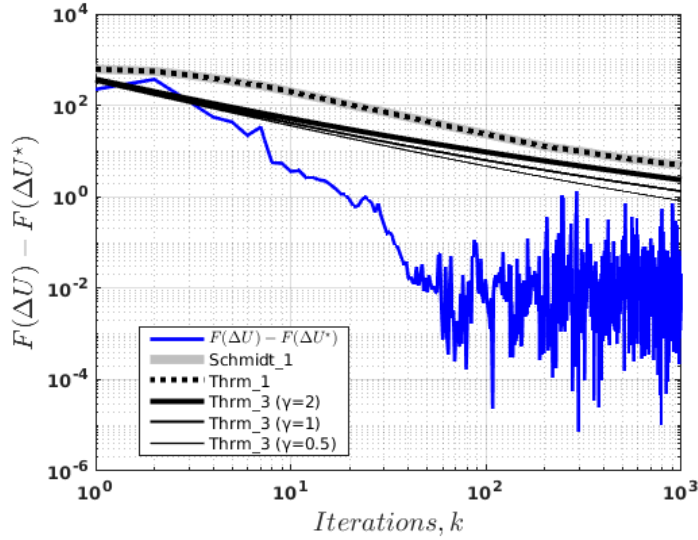


Fig. 1: Upper bounds based on Theorems 1 & 3 vs Proposition 1 ((2.1)) in Schmidt et al. 2010 [25] (with  $\delta = 2.2 \times 10^1; \epsilon_0 = 10^1$ ).

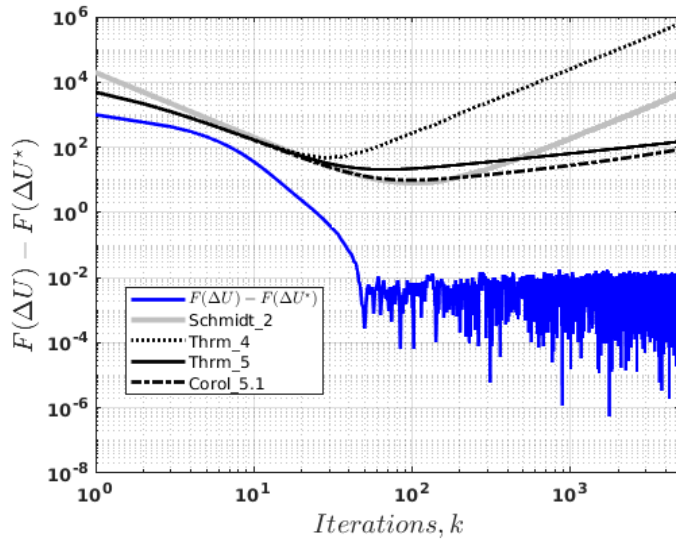


Fig. 2: Upper bounds based on Theorems 4 & 5 vs Proposition 2 in Schmidt et al. 2010 [25] (with  $\delta = 2.2 \times 10^{-4}; \epsilon_0 = 10^{-4}$ ).

691 **Figure 1** suggests that by using the proposed probabilistic bounds, we can predict  
 692 the suboptimality, i.e.,  $f - f^*$ , more accurately and the improvement is more signif-  
 693 icant with lower values of  $\gamma$  (with lower probabilities). Note that the probabilistic  
 694 bounds can possibly drop below the suboptimality plot ( $f - f^*$ ) during some itera-

695 tions; however, this would not present any conflict with the theory as this is what can  
 696 be expected from probabilistic statements (dependent on the parameter  $\gamma$ ) which do  
 697 not hold 100% of the algorithm's execution time.

698 From Figure 2, we can see that none of the bounds can successfully estimate  
 699 the function values suboptimality in the accelerated case, however, the probabilistic  
 700 bound of Theorem 5 gives the best estimate and the slowest divergence rate. The  
 701 bound of Corollary 5.1 slightly improves on Theorem 5 but still diverges, although at  
 702 the slowest rate.

703 **6. Conclusions.** We have analysed the convergence of the proximal gradient de-  
 704 scent under computational errors. We derived deterministic and probabilistic upper  
 705 bounds on the objective function value which we used as an assertion for convergence  
 706 test. We considered the special case in which the gradient  $\nabla g(x^k)$  of  $g$  is computed  
 707 with errors as well as the proximal operator  $\text{prox}_h$  (with respect to  $h$ ) is evaluated  
 708 approximately. We also considered accelerated versions of the proximal gradient de-  
 709 scent, which is known to converge faster in the error-free case, but we have shown that  
 710 this comes at a price of amplified perturbations, which may lead to divergence. We  
 711 proved that the effect of each contributing error term can be decoupled under mild  
 712 assumptions. We also obtained probabilistic bounds with three main advantages:

- 713 • The bounds are sharper (i.e., reflect practical performance better);
- 714 • The bounds are simpler to interpret and predict *a priori*;
- 715 • The contribution of each error term is decoupled.

716 We have also shown that some error terms follow martingale sequences when error  
 717 conditional mean independence and data conditional mean independence assumptions  
 718 both hold. Finally, we have perceived that in the accelerated case, the algorithm actu-  
 719 ally converges to some suboptimal level around the optimum, however, the latter could  
 720 not be determined using the current convergence bounds. This opens the possibility  
 721 of other types of analyses with different error models.

722 **Appendix A. Supplementary results.** The following Lemma establishes  
 723 bounds on the norm of the residual error vector due to proximal error (forward error).

724 LEMMA 1. Consider problem (3.1) and let Assumption 1 hold. For  $L, s > 0$ ,  
 725 define  $G: \mathbb{R}^n \times \mathbb{R}^n \rightarrow (-\infty, \infty]$  as the proper, closed, and  $L$ -strongly convex function

$$726 \quad G(y, x) := g(y) + \nabla g(y)^\top (x - y) + \frac{1}{2s} \|x - y\|_2^2 + h(x),$$

727 Define  $\hat{y}^* := \arg \min G(y, x)$  as the minimizer of  $G$  with respect to  $y$  when  $x$  is fixed,  
 728 and  $y^* \in \{y : G(y, x) - G(\hat{y}^*, x) \leq \epsilon_2\}$  as an  $\epsilon_2$ -approximate solution of the same  
 729 problem. Then,

$$730 \quad \|\hat{y}^* - y^*\|_2 \leq \sqrt{2s\epsilon_2}.$$

731 **THEOREM 6 (Quasi-Fejér monotonicity of the sequence generated by the**  
 732 **proximal gradient method).** Let  $\{x^k\}_{k \geq 0}$  be the sequence generated by the ap-  
 733 proximate proximal gradient (3.7) for solving problem (3.1) under Assumption 1 and  
 734 with  $s_k \leq \frac{1}{L}$ . Assume that, for  $k \geq k_0$ , we have  $\epsilon_2^k \leq c_2 \|x^{k+1} - x^k\|_2 \leq c_2 \rho$  and  
 735  $\|\epsilon_1^k\|_2 \leq c_1 \|\nabla g(x^{k+1}) - \nabla g(x^k)\|_2$ . Then for any  $x^* \in X^*$  and  $k \geq 0$  we have

$$736 \quad (\text{A.1}) \quad \|x^{k+1} - x^*\|_2 \leq \|x^k - x^*\|_2 + \|r^{k+1}\|_2 + s_k \|\epsilon_1^k\|_2 + C_{\rho, 1/L},$$

737 where  $C_{\rho, 1/L} = \sqrt{2Lc_2\rho} + c_1\rho$ . If  $E^{k+1} := \|r^{k+1}\|_2 + s_k \|\epsilon_1^k\|_2$  is a positive and



738 absolutely summable sequence, then  $\{x^k\}_{k \geq 0}$  is a quasi-Féjer sequence relative to the  
 739 set  $X^*$ .

740 *Proof.* we have

(A.2)

$$741 \quad \|x^{k+1} - x^{k_0+1}\|_2 = \left\| \text{prox}_{s_k h}^{\epsilon_2^k}(x^k - s_k \nabla^{\epsilon_1^k} g(x^k)) - \text{prox}_{s_{k_0} h}^{\epsilon_2^{k_0}}(x^{k_0} - s_{k_0} \nabla^{\epsilon_1^{k_0}} g(x^{k_0})) \right\|_2.$$

743 Writing  $\text{prox}_{s_k h}^{\epsilon_2^k}(x)$  as  $\text{prox}_{s_k h}(x) + r^k$  and  $\nabla^{\epsilon_1^k} g(x)$  as  $\nabla g(x) + \epsilon_1^k$  for any suboptimal  
 744 solution  $x^{k_0}$  of (3.1), we obtain

$$745 \quad (A.3) \quad \begin{aligned} \|x^{k+1} - x^{k_0+1}\|_2 = & \left\| \text{prox}_{s_k h}(x^k - s_k \nabla g(x^k) - s_k \epsilon_1^k) \right. \\ & \left. - \text{prox}_{s_{k_0} h}(x^{k_0} - s_{k_0} \nabla g(x^{k_0}) - s_{k_0} \epsilon_1^{k_0}) + r^{k+1} - r^{k_0} \right\|_2. \end{aligned}$$

746 By assumption we have  $\epsilon_2^k \leq c_2 \|x^{k+1} - x^k\|_2$ , or equivalently,

$$747 \quad \|r^{k+1}\|_2 \leq \sqrt{2c_2} \|x^{k+1} - x^k\|_2 / s \text{ and } \epsilon_2^k \leq c_1 \|\nabla g(x^{k+1}) - \nabla g(x^k)\|_2$$

748  $\leq c_1 L \|x^{k+1} - x^k\|_2$  for  $k \geq k_0$ . By the triangle inequality we have

(A.4)

$$749 \quad \begin{aligned} \|x^{k+1} - x^{k_0+1}\|_2 \leq & \left\| \text{prox}_{s_k h}(x^k - s_k \nabla g(x^k) - s_k \epsilon_1^k) \right. \\ & \left. - \text{prox}_{s_{k_0} h}(x^{k_0} - s_{k_0} \nabla g(x^{k_0}) - s_{k_0} \epsilon_1^{k_0}) \right\|_2 + \|r^{k+1}\|_2 + \|r^{k_0+1}\|_2 \\ \leq & \left\| \text{prox}_{s_k h}(x^k - s_k \nabla g(x^k) - s_k \epsilon_1^k) \right. \\ & \left. - \text{prox}_{s_{k_0} h}(x^{k_0} - s_{k_0} \nabla g(x^{k_0}) - s_{k_0} \epsilon_1^{k_0}) \right\|_2 + \|r^{k+1}\|_2 + \sqrt{\frac{2c_2 \rho}{s}} \end{aligned}$$

750 where we have used  $\|x^{k_0+1} - x^{k_0}\|_2 \leq \rho$ .

751 By the nonexpansivity of the proximal operator we have

(A.5)

$$752 \quad \begin{aligned} \|x^{k+1} - x^{k_0+1}\|_2 \leq & \left\| [x^k - s_k \nabla g(x^k)] - [x^{k_0} - s_{k_0} \nabla g(x^{k_0})] \right\|_2 + \|r^{k+1}\|_2 + \sqrt{\frac{2c_2 \rho}{s_{k_0}}} \\ & + s_k \|\epsilon_1^k\|_2 + s_{k_0} \|\epsilon_1^{k_0}\|_2 \\ \leq & \left\| [x^k - s_k \nabla g(x^k)] - [x^{k_0} - s_{k_0} \nabla g(x^{k_0})] \right\|_2 + \|r^{k+1}\|_2 + s_k \|\epsilon_1^k\|_2 \\ & + \sqrt{\frac{2c_2 \rho}{s_{k_0}}} + s_{k_0} c_1 L \rho \end{aligned}$$

753 By the nonexpansivity of the gradient descent operator, i.e.,  $\mathbf{I} - s \nabla g$ , we obtain

$$754 \quad (A.6) \quad \|x^{k+1} - x^{k_0+1}\|_2 \leq \|x^k - x^{k_0}\|_2 + \|r^{k+1}\|_2 + s_k \|\epsilon_1^k\|_2 + C_\rho, \quad \forall s_k \leq \frac{1}{L}$$

$$755 \quad (A.7) \quad = \|x^k - x^{k_0}\|_2 + E^{k+1} + C_\rho,$$

757

$$(A.8) \quad \|x^{k+1} - x^{k_0} - (x^{k_0+1} - x^{k_0})\|_2 \leq \|x^k - x^{k_0}\|_2 + E^{k+1} + C_\rho,$$

760 where  $C_\rho = \sqrt{\frac{2c_2\rho}{s_{k_0}}} + s_{k_0}c_1L\rho$  and  $E^{k+1} = \|r^{k+1}\|_2 + s_k \|\epsilon_1^k\|_2$ .

761 By the triangle difference inequality we have

$$(A.9) \quad \left| \|x^{k+1} - x^{k_0}\|_2 - \|x^{k_0+1} - x^{k_0}\|_2 \right| \leq \|x^k - x^{k_0}\|_2 + E^{k+1} + C_\rho.$$

764 For  $x^{k_0+1} \approx x^{k_0} = x^*$  we have

$$(A.10) \quad \|x^{k+1} - x^*\|_2 \leq \|x^k - x^*\|_2 + E^{k+1} + C_\rho,$$

767 From (A.10) and by [9, Definition 1.1], the sequence  $\{x^k\}_{k \geq 1}$  is quasi-Féjer relative  
768 to the set  $X^*$  if  $\{E^k\}_{k \geq 1}$  is positive and absolutely summable.

769  $\square$

770 **THEOREM 7 (Quasi-Féjer monotonicity of the sequence generated by the**  
771 **accelerated proximal gradient method).** *Let  $\{x^k\}_{k \geq 0}$  be the sequence generated*  
772 *by the approximate accelerated proximal gradient (3.6) for solving problem (3.1) under*  
773 *Assumption 1 and with  $s_k \leq \frac{1}{L}$ . Assume we have summable iterative displacements*  
774  *$\|x^k - x^{k-1}\|_2$  and that, for  $k \geq k_0$ , we have  $\epsilon_2^k \leq c_2 \|x^{k+1} - x^k\|_2 \leq c_2\rho$  and  $\|\epsilon_1^k\|_2 \leq$*   
775  *$c_1 \|\nabla g(x^{k+1}) - \nabla g(x^k)\|_2^\top$ , then for any  $x^{k_0} \in X^{k_0}$  and  $k \geq 0$  we have*

$$(A.11) \quad \|x^{k+1} - x^{k_0}\|_2 \leq \|x^k - x^{k_0}\|_2 + \|x^k - x^{k-1}\|_2 + E^{k+1} + C_{\rho,1/L}$$

777 where  $C_{\rho,1/L} = \sqrt{2Lc_2\rho} + c_1\rho$ ,  $E^{k+1} = \|r^{k+1}\|_2 + s_k \|\epsilon_1^k\|_2$ . If  $E^{k+1} := \|r^{k+1}\|_2 +$   
778  $s_k \|\epsilon_1^k\|_2$  is a positive and absolutely summable sequence, then  $\{x^k\}_{k \geq 0}$  is a quasi-Féjer  
779 sequence relative to the set  $X^{k_0}$ .

780 *Proof.* For any optimal solution  $x^{k_0}$  of (3.1), we have

$$(A.12)$$

$$\|x^{k+1} - x^{k_0+1}\|_2 = \left\| \text{prox}_{s_k h}^{\epsilon_2^k}(y^k - s_k \nabla^{\epsilon_1^k} g(y^k)) - \text{prox}_{s_{k_0} h}^{\epsilon_2^{k_0}}(x^{k_0} - s_{k_0} \nabla^{\epsilon_1^{k_0}} g(x^{k_0})) \right\|_2.$$

783 Rewriting  $\text{prox}_{s_k h}^{\epsilon_2^k}(y)$  as  $\text{prox}_{s_k h}(y) + r^k$  and  $\nabla^{\epsilon_1^k} g(y)$  as  $\nabla g(y) + \epsilon_1^k$  we obtain

$$(A.13) \quad \begin{aligned} \|x^{k+1} - x^{k_0+1}\|_2 &= \left\| \text{prox}_{s_k h}(y^k - s_k \nabla g(y^k) - s_k \epsilon_1^k) \right. \\ &\quad \left. - \text{prox}_{s_{k_0} h}(x^{k_0} - s_{k_0} \nabla g(x^{k_0}) - s_{k_0} \epsilon_1^{k_0}) + r^{k+1} - r^{k_0} \right\|_2. \end{aligned}$$

785 By assumption we have  $\epsilon_2^k \leq c_2 \|x^{k+1} - x^k\|_2$  and

786  $\epsilon_2^k \leq c_1 \|\nabla g(x^{k+1}) - \nabla g(x^k)\|_2 \leq c_1 L \|x^{k+1} - x^k\|_2$  for  $k \geq k_0$ . By the triangle

787 inequality we have

(A.14)

$$\begin{aligned} \|x^{k+1} - x^{k_0+1}\|_2 &\leq \left\| \text{prox}_{s_k h}(y^k - s_k \nabla g(y^k) - s_k \epsilon_1^k) \right. \\ &\quad \left. - \text{prox}_{s_{k_0} h}(x^{k_0} - s_{k_0} \nabla g(x^{k_0}) - s_{k_0} \epsilon_1^{k_0}) \right\|_2 + \|r^{k+1}\|_2 + \|r^{k_0+1}\|_2 \\ &\leq \left\| \text{prox}_{s_k h}(y^k - s_k \nabla g(y^k) - s_k \epsilon_1^k) \right. \\ &\quad \left. - \text{prox}_{s_{k_0} h}(x^{k_0} - s_{k_0} \nabla g(x^{k_0}) - s_{k_0} \epsilon_1^{k_0}) \right\|_2 + \|r^{k+1}\|_2 + \sqrt{\frac{2c_2\rho}{s}} \end{aligned}$$

789 where we have used  $\|x^{k_0+1} - x^{k_0}\|_2 \leq \rho$ .

790 By the nonexpansivity of the proximal operator we have

$$\begin{aligned} \|x^{k+1} - x^{k_0+1}\|_2 &\leq \| [y^k - s_k \nabla g(y^k)] - [x^{k_0} - s_{k_0} \nabla g(x^{k_0})] \|_2 + \|r^{k+1}\|_2 \\ &\quad + \sqrt{\frac{2c_2\rho}{s_{k_0}}} + s_k \|\epsilon_1^k\|_2 + s_{k_0} \|\epsilon_1^{k_0}\|_2 \\ &\leq \| [y^k - s_k \nabla g(y^k)] - [x^{k_0} - s_{k_0} \nabla g(x^{k_0})] \|_2 + \|r^{k+1}\|_2 \\ &\quad + s_k \|\epsilon_1^k\|_2 + \sqrt{\frac{2c_2\rho}{s_{k_0}}} + s_{k_0} c_1 L \rho \end{aligned}$$

792 By the nonexpansivity of the gradient descent operator, i.e.,  $\mathbf{I} - s\nabla g$ , we obtain

$$\begin{aligned} \|x^{k+1} - x^{k_0+1}\|_2 &\leq \|y^k - x^{k_0}\|_2 + \|r^{k+1}\|_2 + s_k \|\epsilon_1^k\|_2 + C_{\rho, s_{k_0}}, \quad \forall s_k \leq \frac{1}{L} \\ &= \|x^k - x^{k_0} + \beta_k(x^k - x^{k-1})\|_2 + E^{k+1} + C_{\rho, s_{k_0}} \\ &= \|x^k - x^{k_0}\|_2 + \|x^k - x^{k-1}\|_2 + E^{k+1} + C_{\rho, s_{k_0}}, \end{aligned}$$

794 where  $C_{\rho, s_{k_0}} = \sqrt{\frac{2c_2\rho}{s_{k_0}}} + s_{k_0} c_1 L \rho$ ,  $E^{k+1} = \|r^{k+1}\|_2 + s_k \|\epsilon_1^k\|_2$  and we used  $\beta_k \leq 1$ .

795 By the triangle difference inequality we have

(A.17)

$$\| \|x^{k+1} - x^{k_0}\|_2 - \|x^{k_0+1} - x^{k_0}\|_2 \| \leq \|x^k - x^{k_0}\|_2 + \|x^k - x^{k-1}\|_2 + E^{k+1} + C_{\rho, s_{k_0}},$$

797 For  $x^{k_0+1} \approx x^{k_0} = x^*$  we have

(A.18)

$$\|x^{k+1} - x^*\|_2 \leq \|x^k - x^*\|_2 + \|x^k - x^{k-1}\|_2 + E^{k+1} + C_{\rho, s_{k_0}},$$

799 From (A.18) and by [9, Definition 1.1], the sequence  $\{x^k\}_{k \geq 1}$  is quasi-Féjer relative  
800 to the set  $X^*$  if  $\{E^k\}_{k \geq 1}$  is positive and absolutely summable provided we have  
801 summable iterative displacements  $\|x^k - x^{k-1}\|_2$ .  $\square$

802 **Acknowledgements.** This work was supported by the Engineering and Physi-  
803 cal Research Council (EPSRC) grants (EP/T026111/1 and EP/S000631/1), and the  
804 MOD University Defence Research Collaboration (UDRC).

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